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(54) METHOD AND DEVICE FOR CODING THE LOW DENSITY GENERATOR MATRIX CODE

VERFAHREN UND EINRICHTUNG ZUM CODIEREN DES LOW-DENSITY-GENERATOR-MATRIXCODES

PROCÉDÉ ET DISPOSITIF POUR CODER LE CODE DE MATRICE DE GÉNÉRATEUR À FAIBLE DENSITÉ

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Description

TECHNICAL FIELD

[0001] The present invention relates to the field of communications, especially to a method and device for encoding the low density generator matrix code.

BACKGROUND

[0002] Erasure channel is an important type of channel model, for example, files are transmitted on the Internet based on data packet communications, and generally each data packet is either received by a receiver correctly or not received by the receiver at all. Error detection and retransmission mechanism adopted in Transmission Control Protocol (TCP) regarding to network packet loss is to control the data packets that need to be retransmitted by using the feedback channel from an input end to an output end. When the receiver detects packet loss, it will generate a retransmitting control signal, until the data packets are received completely and correctly; when the receiver receives a data packet, it will likewise generate a receiving confirmation signal. Also the sender will track every data packet until a notification signal is fed back, otherwise it will resend the data packet.

[0003] A data broadcasting service based on a stream mode and a file download mode is a point-to-multipoint service, in which feedback is not allowed, so traditional error detection and retransmission mechanism may not be applied, and Forward Error Correction (FEC) has to be adopted to ensure reliable transmission. FEC codes of the classical application layer include Reed-Solomon (RS) codes, digital Fountain codes and so on. Encoding/decoding of RS codes is of high complexity, so RS codes are usually applied to short codes. Luby Transform (LT) codes and Raptor codes are two types of digital Fountain codes that have practical applications. LT codes have linear encoding and decoding time, with essential enhancement in comparison with RS codes, but Raptor codes adopts precoding technique, so Raptor codes have a greater decoding efficiency. Raptor codes are adopted as a FEC encoding scheme in both 3GPP Multimedia Broadcast/Multicast Service (MBMS) and Digital Video Broadcasting (DVB).

[0004] As a group of code blocks with constant length, linear block codes are denoted as (n, k) and are usually used for FEC. During the encoding, k information bits are encoded into an n-bit code block. Since 2^k code words of the block codes (n, k) constitute a k-dimensional subspace, the 2^k code words may definitely be generated by k linearly independent basements, and if the k linearly independent basements are written in a form of matrix, then

$$G = \begin{bmatrix} g_{1,n-1} & g_{1,n-2} & \cdots & g_{1,0} \\ g_{2,n-1} & g_{2,n-2} & \cdots & g_{2,0} \\ \vdots & \vdots & & \vdots \\ g_{k,n-1} & g_{k,n-2} & \cdots & g_{k,0} \end{bmatrix},$$

wherein, any code word in the block codes (n, k) may be generated by a linear combination of this group of basements, namely

$$C = m.G = [m_1 \quad m_2 \quad \ldots \quad m_{n-k}] \cdot \begin{bmatrix} g_{1,n-1} & g_{1,n-2} & \cdots & g_{1,0} \\ g_{2,n-1} & g_{2,n-2} & \cdots & g_{2,0} \\ \vdots & \vdots & & \vdots \\ g_{k,n-1} & g_{k,n-2} & \cdots & g_{k,0} \end{bmatrix}.$$

[0005] Here G is called a generator matrix of codes. Obviously, for each row of the generator matrix only linear independence needs to be satisfied (disregarding of the minimum distance), but the basements of a k-dimensional space may be composed of k selected linearly independent vectors, so the generator matrix G of codes is not unique, no matter which form is adopted, however, they generate the same subspace, namely the same (n, k) block codes.

[0006] If the first k bits of the encoded code words are the same as information bits, this code is called systematic code. The encoding procedure is that of generating n-bit code length by k information bits, and error detection and error correction are implemented by adding n-k check bits.

[0007] LT codes are also encoded by using the sparsity of the generator matrix, in comparison with the Low Density

Generator Matrix Code (LDGC) of ZTE Corporation, LT codes do not support encoding method of systematic codes, so it is difficult for LT codes to meet the requirement of some practical FEC encoding; Raptor codes support systematic codes, but need a separate precoding procedure, namely need a precoding matrix, therefore, leading to a higher encoding complexity, while the LDGC directly uses a generator matrix to encode without needing an extra precoding matrix, and the LDGC takes use of back substitution method to solve an upper (or a lower) triangular function while encoding, so the encoding complexity of the LDGC is much lower than that of Raptor codes. In a word, compared with LT codes, the LDGC has an advantage in supporting systematic codes; compared with Raptor codes, the LDGC has an advantage in lower encoding complexity.

[0008] In the reference document "L. H. C. Lee, 'Error-Control Block Codes for Communications Engineers, Linear Block Codes', XP002642221", it deals with the theory of error-control block codes. The treatment here concentrates on the basic principles of encoding and decoding binary linear block codes.

[0009] In the reference document "Vincent Roca, Christoph Neumann: 'Design, Evaluation and Comparison of Four Large Block FEC Codes, LDPC, LDGM, LDGM Staircase and LDGM Triangle, plus a Reed-Solomon Small Block FEC Codec', internet article, 9 June 2004 (2004-06-09), Retrieved from the Internet: URL: http://hal. Inria. Fr/inria-00070770/document", it introduces with a great level of details the design and the performance evaluation of four large block FEC codes: regular LDPC (Low Density Parity Check), LDGM (Low Density Generator Matrix), LDGM Staircase and LDGM Triangle, that are all capable of operating on source blocks that are composed of several tens of thousand of packets. If all four belong to the category of codes on graphs, they differ by their encoding algorithm. We evaluate and compare the performances of all four codes, as well as a Reed-Solomon codec which serves as a reference and gives an idea of the level of performance achievable with small block codes. The comparison focuses on three major aspects: the global decoding inefficiency (caused either by the FEC code or by the need to split the object into several blocks), the encoding and decoding speeds, and the memory requirements at the encoder and decoder. From these performance results, we identify the situation where they best operate.

SUMMARY

[0010] In order to reduce the encoding complexity and obtain a better code-performance, a method and device for encoding the low density generator matrix code are provided in the present invention.

[0011] According to one aspect of the present invention, a method for encoding the low density generator matrix code is provided. The method for encoding the low density generator matrix code according to embodiments of the present invention comprises the following steps: constructing a generator matrix G_{ldgc} with L rows and N+L-K columns, wherein, a phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper triangular matrix or a lower triangular matrix, the K, L and N being positive integers, and K<L<N; adding L-K known bits to a K-length information bit sequence that needs to be encoded, generating an information bit sequence m with a length of L; according to $I \times G_{ldgc}(1:L, 1:L)=m$, using the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} and the L-length information bit sequence m, generating an intermediate variable I by means of upper/lower triangular function solver, and according to $C=I \times G_{ldgc}$, encoding the intermediate variable I by using the generator matrix G_{ldgc}, and generating an encoding sequence that has a length of N+L-K; and deleting L-K known bits from the encoding sequence that has a length of N+L-K, generating an encoding sequence that has a length of N.

[0012] Wherein, when the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-left triangular matrix or a lower-left triangular matrix, the L-K known bits are added in front of the K-length information bit sequence that is to be encoded. When the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-right triangular matrix or a lower-right triangular matrix, the L-K known bits are added after the K-length information bit sequence that is to be encoded. Wherein, a column weight of the generator matrix G_{ldgc} is in accordance with a degree distribution principle similar to that of LT codes.

[0013] According to another aspect of the present invention, a method for encoding the low density generator matrix code is provided. The method for encoding the low density generator matrix code according to embodiments of the present invention comprises the following steps: constructing a generator matrix G_{ldgc} with K rows and N columns, wherein, a phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} is an upper triangular matrix or a lower triangular matrix, the K and N being positive integers, and K<N; according to $I \times G_{ldgc}(1:K, 1:K) =s$, using a K-length information bit sequence s that needs to be encoded and the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc}, generating an intermediate variable I by means of upper/lower triangular function solver; according to $C=I \times G_{ldgc}$, encoding the intermediate variable I by using the generator matrix G_{ldgc}, and generating a code word that has a length of N. Wherein, a column weight of the generator matrix G_{ldgc} is in accordance with a degree distribution principle similar to that of LT codes. The phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} may be an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix.

[0014] According to another aspect of the present invention, a device for encoding the low density generator matrix

code is provided. The device for encoding the low density generator matrix code according to embodiments of the present invention comprises: a matrix generating unit, being used for generating a generator matrix G_{ldgc} with L rows and N+L-K columns, outputting the generated matrix G_{ldgc} to a block code encoding unit, and outputting a phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} to a precoding unit, wherein, the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper triangular matrix or a lower triangular matrix, the K, L and N being positive integers, and $K<L<N$; a bit filling unit, being used for adding L-K known bits to a K-length information bit sequence that is to be encoded, generating an information bit sequence m with a length of L, and outputting the L-length information bit sequence m to the precoding unit; the precoding unit, being used for according to $I \times G_{ldgc}(1:L, 1:L)=m$, using the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} and the L-length information bit sequence m, generating an intermediate variable I by means of upper/lower triangular function solver, and outputting the intermediate variable I to the block code encoding unit (the precoding unit may be a unit for solving upper and/or lower triangular function, with two inputs, one being $G_{ldgc}(1:L, 1:L)$ and the other being m, and one output that is I); the block code encoding unit, being used for according to $C=I \times G_{ldgc}$, using the generator matrix G_{ldgc} to encode the intermediate variable I, generating an encoding sequence that has a length of N+L-K (the block code encoding unit may be a matrix square unit, with two inputs, which are respectively G_{ldgc} and I, and one output that is C); and a bit deleting unit, being used for deleting L-K known bits from the encoding sequence that has a length of N+L-K, generating an encoding sequence that has a length of N.

[0015] Wherein, the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} may be an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix or a lower-right triangular matrix. When the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-left triangular matrix or lower-left triangular matrix, the bit filling unit adds L-K known bits in front of the K-length information bit sequence that needs to be encoded. When the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-right triangular matrix or lower-right triangular matrix, the bit filling unit adds L-K known bits after the K-length information bit sequence that needs to be encoded. A column weight of the generator matrix G_{ldgc} is in accordance with a degree distribution principle similar to that of LT codes.

[0016] According to another aspect of the present invention, a device for encoding the low density generator matrix code is provided. The device for encoding the low density generator matrix code according to embodiments of the present invention comprises: a matrix generating unit, being used for generating a generator matrix G_{ldgc} with K rows and N columns, wherein, a phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} is an upper triangular matrix or a lower triangular matrix, the K and N being positive integers, and $K<N$; a precoding unit, being used for according to $I \times G_{ldgc}(1:K, 1:K)=s$, using a K-length information bit sequence s that is to be encoded and the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc}, generating an intermediate variable I by means of upper/lower triangular function solver; and a block encoding unit, being used for according to $C=I \times G_{ldgc}$, using the generator matrix G_{ldgc} to encode the intermediate variable I, and generating a code word that has a length of N.

[0017] Wherein, the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} may be an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix. The column weight of generator matrix G_{ldgc} is in accordance with the degree distribution principle similar to that of LT codes.

[0018] According to the present invention, a better code-performance can be obtained while reducing encoding complexity.

BRIEF DESCRIPTION OF THE DRAWINGS

[0019] The drawings illustrated hereby are used for providing further understanding of the present invention, as one part of this application, and the exemplary embodiments and illustration thereof are used to explain the present invention, but not to confine the present invention improperly. In the drawings:

Fig. 1 is a flowchart illustrating the LDGC encoding method according to an embodiment of the present invention;

Fig. 2 is a block diagram illustrating the LDGC encoding device according to an embodiment of the present invention;

Fig. 3 is a schematic diagram illustrating the LDGC generator matrix;

Fig. 4 is an example of the LDGC generator matrix with filled bits according to an embodiment of the present invention;

Fig. 5 is a flowchart illustrating the LDGC encoding method according to another embodiment of the present invention;

Fig. 6 is a block diagram illustrating the LDGC encoding device according to another embodiment of the present invention;

Fig. 7 is an example of the LDGC generator matrix without filled bits according to another embodiment of the present invention;

Fig. 8 is a schematic diagram illustrating the precoding unit in the LDGC encoding device according to another embodiment of the present invention;

Fig. 9 is a schematic diagram illustrating the block code encoding unit in the LDGC encoding device according to another embodiment of the present invention.

DETAILED DESCRIPTION

[0020] The Low Density Generator Matrix Code (LDGC) is a kind of linear block codes, and non-zero elements of a LDGC generator matrix are usually sparse. Meanwhile, the LDGC is also a kind of systematic code, and a phalanx composed by the first k columns of the LDGC generator matrix is usually an upper triangular matrix or a lower triangular matrix, inversion of the matrix may be obtained by the method of iteration. The encoding procedure of LDGC includes getting an intermediate variable by using a corresponding relation between an information bit of systematic code and the intermediate variable firstly, and then multiplying the intermediate variable by the generator matrix to get encoded code words. The decoding procedure of LDGC includes getting the intermediate variable by using the generator matrix firstly, and then figuring out the information bit according to the transformation relation between the information bit and the intermediate variable. Embodiments of the present invention will be illustrated in detail with reference to the drawings.

[0021] With reference to Fig. 1, which illustrates the procedure of encoding a K-length information bit sequence, and then outputting an N-length encoded code word bit sequence to a subsequent processing unit. Wherein, the length of the check bit is M=N-K, and the code rate is r=K/N. As shown in Fig. 1, the LDGC encoding method according to an embodiment of the present invention comprises the following steps:

S102, a generator matrix G_{ldgc} with L rows and N+L-k columns is constructed. Wherein, a phalanx $G_{ldgc}(1:L, 1:L)$ that consists of the first L columns and all rows of this generator matrix is an upper triangular matrix (or a lower triangular matrix), the K, L and N being given positive integers, and K<L<N;

S104, d=L-K known filling bits are added to a 1*K information bit sequence s, to construct a 1*L information bit sequence m;

S106, LDGC is systematic code, so $I \times G_{ldgc}(1:L, 1:L)=m$. Since the $G_{ldgc}(1:L, 1:L)$ is an upper triangular matrix (or a lower triangular matrix), an intermediate variable I of 1*K is obtained by solving the function, and according to $C=I \times G_{ldgc}$, the intermediate variable I is encoded to get a 1*(N+d) code word;

S108, the filling bits added in the step S104 are deleted from the 1*(N+d) code word generated in the above step, and finally an N-bit encoded code word is obtained and sent out.

[0022] Wherein, in each column of the generator matrix, a number of element "1" (i.e. a column weight) must be in accordance with a certain degree distribution principle. The phalanx composed by the first L columns and all rows of the generator matrix may be an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix (as shown in Fig. 3).

[0023] Wherein, if $G_{ldgc}(1:L, 1:L)$ is an upper-left triangular matrix or a lower-left triangular matrix, then for the information bit sequence whose initial length is K,, d=L-K known filling bits are added in front of the K information bit sequences. If $G_{ldgc}(1:L, 1:L)$ is an upper-right triangular matrix or a lower-right triangular matrix, then for the information bit sequence whose initial length is K, d=L-K known filling bits are added after the K information bit sequences. To be noted, positions where the filling bits are added are not limited as above.

[0024] With reference to Fig. 2, an LDGC encoding device according to an embodiment of the present invention is illustrated. The LDGC encoding device is used to encode an input K-bit binary information bit stream, and then output an N-bit binary code word bit stream to the subsequent processing unit. As shown in Fig. 2, the LDGC encoding device comprises:

a matrix generating unit 202, being used for generating a generator matrix G_{ldgc} with L rows and N+L-K columns. Wherein, a phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L columns and all rows of the generator matrix G_{ldgc} is an upper

triangular matrix (or a lower triangular matrix). The matrix generator outputs $G_{ldgc}(1:L, 1:L)$ to a subsequent precoding unit, and outputs $G_{ldgc}(1:L, 1:N+L-K)$ to a subsequent block code encoding unit. Wherein, the K, L and N are given positive integers, and K<L<N;

a bit filling unit 204, being used for adding d=L-K known filling bits to an input 1*K information bit sequence s, generating a 1*L information bit stream m, and outputting the generated information bit stream m to the precoding unit;

the precoding unit 206, being used for generating a 1*L intermediate variable I by solving the function of the 1*L information bit stream m, and outputting an intermediate variable I to a block code encoder (as shown in Fig. 8, the precoding unit may be a unit for solving an upper and/or a lower triangular function, with two inputs, one being $G_{ldgc}(1:L, 1:L)$, and the other being the information bit sequence m that has a length of L, and one output that is the intermediate variable I);

the block code encoding unit 208, being used for encoding the intermediate variable I, generating a 1*(N+d) binary code word bit stream C, and outputting C to a bit deleting unit (as shown in Fig. 9, the block code encoding unit may be a unit for the matrix multiplying, with two inputs, one being G_{ldgc} and the other being the L-length intermediate variable I, and the output is the encoded code word C);

a bit deleting unit 210, being used for deleting the filled d filling bits from the bit filling unit, and finally obtaining an N-bit encoded code word.

[0025] Wherein, the matrix generating unit determines a generator matrix according to the following rules: in every column of the generator matrix, the number of element "1" (i.e. the column weight) must be in accordance with a certain degree distribution principle. The phalanx composed by the first L columns and all rows of the generator matrix is an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix (as shown in Fig. 3).

[0026] Wherein, when $G_{ldgc}(1:L, 1:L)$ is an upper-left triangular matrix or a lower-left triangular matrix, the bit filling unit adds d=L-K known filling bits in front of the K information bit sequences. When $G_{ldgc}(1:L, 1:L)$ is an upper-right triangular matrix or a lower-right triangular matrix, the bit filling unit adds d=L-K known filling bits after the K information bit sequences. To be noted, positions where the filling bits are added are not limited as above.

[0027] Wherein, because LDGC is a systematic code with $I \times G_{ldgc}(1:L, 1:L)=m$, the precoding unit solves the function when $G_{ldgc}(1:L, 1:L)$ is an upper triangular matrix (or a lower triangular matrix) to obtain the 1*L intermediate variable I.

[0028] Wherein, the block code encoding unit encodes the intermediate variable I according to $C=I \times G_{ldgc}$, and gets a 1*(N-d) binary code word stream.

[0029] For example, there is a 1*K=1*24 binary information bit data stream s (s may be denoted as D8AB13 in hexadecimal format) that is to be encoded by the encoding device according to the embodiments of the present invention into 72-bit LDGC code words, so K=24, N=72.

[0030] The matrix generating unit generates a generator matrix G_{ldgc} with L=48 rows and N+L-K=96 columns. Wherein, the G_{ldgc} adopts a degree distribution principle similar to that of LT codes, and a phalanx $G_{ldgc}(1:L, 1:L)=G_{ldgc}(1:48, 1:48)$ composed by the first L columns and all rows of the G_{ldgc} is an upper-right triangular matrix, as shown in Fig. 4 (wherein, black spots denote element "1", while blank positions denote element "0"). The matrix generating unit outputs $G_{ldgc}(1:L, 1:L)=G_{ldgc}(1:48, 1:48)$ to the subsequent precoding unit, and outputs $G_{ldgc}(1:L, 1:N+L-K)= G_{ldgc}(1:48, 1:96)$ to the subsequent block code encoding unit.

[0031] The bit filling unit adds d=L-K=24 known filling bits p (p is denoted as 9A0C2C in hexadecimal format) after the 1*K=1*24 input information bit stream, to generate a 1*L=1*48 information bit stream m (m is denoted as D8AB139A0C2C in hexadecimal format), and outputs m to the precoding unit.

[0032] According to the characteristic of the LDGC that is a systematic code ($I \times G_{ldgc}(1:48, 1:48)=m$, because the LDGC is a systematic code), and $G_{ldgc}(1:L, 1:L)=G_{ldgc}(1:48, 1:48)$ is an upper-right triangular matrix, the precoding unit solves the function of the input 1*K=1*24 information bit stream s to get a 1*L=1*48 intermediate variable I (I is denoted as 942DA94E0A24 in hexadecimal format), and outputs I to the block code encoding unit.

[0033] According to $C=I \times G_{ldgc}$, the block code encoding unit encodes the input intermediate variable I to generate a 1*(N+d)=1*96 binary code word c (c is denoted as D8AB139A0C2CCD3AC516ED52 in hexadecimal format), and outputs c to the bit deleting unit.

[0034] The bit deleting unit deletes the added d=24 filling bits from the 1*(N+d)=1*96 binary code word c, and finally gets an N=72 bits encoded code word (denoted as D8AB13CD3AC516ED52 in hexadecimal format) and sends it out.

[0035] In order to reduce the complexity of encoding, another LDGC encoding method is provided in the present invention. The encoding method is different from the above encoding method in that the bit filling and bit deleting are not comprised. With reference to Fig. 5, the procedure of encoding a K-length information bit sequence and then outputting

an encoded code word bit sequence with a length of N to the subsequent processing unit is illustrated. Wherein, the length of check bit is M=N-K, and code rate is r=K/N. As shown in Fig. 5, the LDGC encoding method according to another embodiment of the present invention comprises the following steps:

S502, a generator matrix G_{ldgc} with K rows and N columns is constructed. Wherein, a phalanx $G_{ldgc}(1:K, 1:K)$ consisting of all rows and the first K columns of the generator matrix G_{ldgc} is an upper triangular matrix (or a lower triangular matrix), the K and N being given positive integers, and K<N;

S504, as LDGC is a systematic code, then $I \times G_{ldgc}(1:K, 1:K)=s$, since the phalanx $G_{ldgc}(1:K, 1:K)$ is an upper triangular matrix (or a lower triangular matrix), an intermediate variable I is obtained by solving the function;

S506, according to $C=I \times G_{ldgc}$, the intermediate variable I is encoded to get a 1*N code word.

[0036] Wherein, in each column of the generator matrix, the number of element "1" (i.e. the column weight) must be in accordance with a certain degree distribution principle; the phalanx composed by the first K columns and all rows of the generator matrix is an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix (as shown in Fig. 3).

[0037] With reference to Fig. 6, a LDGC encoding device according to another embodiment of the present invention is illustrated. The encoding device is used for encoding an input K-bit binary information bit stream, and then outputting an N-bit binary code word bit stream to the subsequent processing unit. As shown in Fig. 6, the LDGC encoding device comprises:

a matrix generating unit 602, being used for generating a generator matrix G_{ldgc} with K rows and N columns. Wherein, a phalanx $G_{ldgc}(1:K, 1:K)$ that consists of the first K columns and all rows of the generator matrix G_{ldgc} is an upper triangular matrix (or a lower triangular matrix). The matrix generating unit outputs $G_{ldgc}(1:K, 1:K)$ to a subsequent precoding unit, and outputs $G_{ldgc}(1:K, 1:N)$ to a subsequent block code encoding unit. Wherein, K<N;

the precoding unit 604, being used for generating an input 1*K intermediate variable I by solving a function of a 1*K information bit stream s, and outputting an intermediate variable I to the block code encoding unit (as shown in Fig. 8, the precoding unit may be a unit for solving an upper and/or a lower triangular function, with two inputs, one being $G_{ldgc}(1:L, 1:L)$ and the other being an L-length information bit sequence m);

the block code encoding unit 606, being used for encoding the input intermediate variable I, generating a 1*N binary code word C and sending C out (as shown in Fig. 9, the block code encoding unit may be a unit for the matrix multiplying, with two inputs, one being G_{ldgc}, and the other being the L-length intermediate variable I).

[0038] Wherein, the matrix generating unit determines a generator matrix according to the following rules: in every column of the generator matrix, the number of element "1" (i.e. the column weight) must be in accordance with a certain degree distribution principle. The phalanx composed by the first K columns and all rows of the generator matrix is an upper-left triangular matrix, a lower-left triangular matrix, an upper-right triangular matrix, or a lower-right triangular matrix (as shown in Fig. 3).

[0039] Wherein, according to that the LDGC is a systematic code (as the LDGC is a systematic code, then $I \times G_{ldgc}(1:K, 1:K)=s$), the precoding unit solves the function when $G_{ldgc}(1:K, 1:K)$ is an upper triangular matrix (or a lower triangular matrix) to obtain a 1*K intermediate variable I, and outputs I.

[0040] Wherein, the block code encoding unit encodes the intermediate variable I according to $C=I \times G_{ldgc}$, gets a 1*N binary code word stream and outputs it.

[0041] For example, there is a 1*K=1*24 binary information bit data stream s (s may be denoted as D99274 in hexadecimal format) that is to be encoded into 72-bit LDGC by the encoding device according to another embodiment of the present invention, wherein K=24, N=72.

[0042] The matrix generating unit generates a generator matrix G_{ldgc} with K=24 rows and N=72 columns. Therein, the G_{ldgc} adopts a degree distribution principle similar to that of LT codes, and a phalanx $G_{ldgc}(1:K, 1:K)=G_{ldgc}(1:24, 1:24)$ composed by the first K columns and all rows of G_{ldgc} is an upper-right triangular matrix, as shown in Fig. 7 (therein, the black spots denote element "1", while blank positions denote element "0"). The matrix generating unit outputs $G_{ldgc}(1:K, 1:K)=G_{ldgc}(1:24, 1:24)$ to the subsequent precoding unit, and outputs $G_{ldgc}(1:K, 1:N)= G_{ldgc}(1:24, 1:72)$ to the subsequent block code encoding unit.

[0043] According to that the LDGC is a systematic code and $I \times G_{ldgc}(1:L, 1:L)=m$, and since $G_{ldgc}(1:K, 1:K)=G_{ldgc}(1:24, 1:24)$ is an upper-right triangular matrix, the precoding unit solves the function of the input 1*K information bit stream s to generate a 1*K=1*24 intermediate variable I (I may be denoted as B4B304 in hexadecimal format), and outputs I to

the block code encoding unit.

[0044] According to $C=I\times G_{ldgc}$, the block encoding unit encodes the input intermediate variable I to generate a 1*N=1*72 binary code word c (c is denoted as D99274A593CC1AC461 in hexadecimal format), and sends c out.

[0045] To be noted, although $G_{ldgc}(1:48,1:48)$ and $G_{ldgc}(1:24,1:24)$ are illustrated to be the upper-right triangular matrixes, other forms of matrixes are also applicable for the present invention. Accordingly, positions where filling bits are added are not limited as above. If only an intermediate vector is obtained using the characteristic of the systematic code, and the codes are generated and output via the intermediate vector, the same technical effects may be obtained.

[0046] Besides, still to be noted, although the degree distribution of the generator matrix G_{ldgc} is illustrated to be similar to that of LT codes in the above embodiments, the present invention is not limited as above, and the same technical effects may be obtained if only the LDGC matrix is generated according to a certain degree distribution.

[0047] In sum, the present invention may support encoding of any information block length and any code rate, and is similar to Raptor codes in performance, both of which can reach the optimal theoretical performance.

[0048] The present invention is not limited to the above described exemplary embodiments. For those skilled in the art, various modifications and changes may be made to the present invention. Any modification, equivalent substitute and improvement within the scope of the present invention is deemed to be included within the scope of the present invention, as defined by the claims.

Claims

1. A method for encoding the Low Density Generator Matrix Code, LDGC, comprising the following steps:

   constructing (S102) a generator matrix G_{ldgc} with L rows and N+L-K columns, wherein, an phalanx $G_{ldgc}(1:L, 1:L)$ consisting of L rows and the first L columns of the generator matrix G_{ldgc} being an upper triangular matrix or a lower triangular matrix, the K, L and N being positive integers, and K<L<N;
   adding (S104) L-K known bits to a K-length information bit sequence that is to be encoded, and generating an information bit sequence m with a length of L;
   according to $I\times G_{ldgc}(1:L, 1:L)=m$, using the phalanx $G_{ldgc}(1:L, 1:L)$ consisting of L rows and the first L columns of the generator matrix G_{ldgc} and the information bit sequence m with a length of L, generating an intermediate variable I by means of upper/lower triangular function solver, and according to $C = I\times G_{ldgc}$, using the generator matrix G_{ldgc} to encode the intermediate variable I, and generating an encoding sequence with a length of N+L-K (S106); and
   deleting L-K known bits from the encoding sequence with a length of N+L-K, and generating an encoding sequence with a length of N (S108).

2. The encoding method according to claim 1, wherein, when the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-left triangular matrix or a lower-left triangular matrix, adding the L-K known bits in front of the K-length information bit sequence that is to be encoded.

3. The encoding method according to claim 1, wherein, when the phalanx $G_{ldgc}(1:L, 1:L)$ which consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-right triangular matrix or a lower-right triangular matrix, adding the L-K known bits after the K-length information bit sequence that is to be encoded.

4. The encoding method according to any of claims 1 to 3, wherein, a column weight of the generator matrix G_{ldgc} being in accordance with a degree distribution principle similar to that of LT codes.

5. A method for the encoding Low Density Generator Matrix Code, LDGC, comprising the following steps:

   Constructing (S502) a generator matrix G_{ldgc} with K rows and N columns, wherein a phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} being an upper triangular matrix or a lower triangular matrix, the K and N being positive integers, and K<N;
   according to $I\times G_{ldgc}(1:K, 1:K)=s$, using the K-length information bit sequence s that is to be encoded and the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc}, generating (S504) an intermediate variable I by means of upper/lower triangular function solver; and
   according to $C = I\times G_{ldgc}$, encoding the intermediate variable I by using the generator matrix G_{ldgc}, and generating a code word with a length of N (S506).

6. The encoding method according to claim 5, wherein, the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the

first K columns of the generator matrix G_{ldgc} being an upper-left triangular matrix, or a lower-left triangular matrix, or an upper-right triangular matrix, or a lower-right triangular matrix.

7. The encoding method according to claim 5 or claim 6, wherein, a column weight of the generator matrix G_{ldgc} being in accordance with a degree distribution principle similar to that of LT codes.

8. A device for the encoding Low Density Generator Matrix Code, LDGC, comprising:

a matrix generating unit (202), being used for generating a generator matrix G_{ldgc} with L rows and N+L-K columns, outputting the generated matrix G_{ldgc} to a block code encoding unit (208), and outputting a phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} to a precoding unit (206), wherein, the phalanx $G_{ldgc}(1:L, 1:L)$ that consists L rows and the first L columns of the generator matrix G_{ldgc} being an upper triangular matrix or a lower triangular matrix, the K, L and N being positive integers, and $K<L<N$;
a bit filling unit (204), being used for adding L-K known bits to a K-length information bit sequence that is to be encoded, generating an information bit sequence m with a length of L, and outputting the generated information bit sequence m with a length of L to the precoding unit (206);
the precoding unit (206), being used for according to $I \times G_{ldgc}(1:L, 1:L)=m$, by using the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} and the information bit sequence m with a length of L, generating an intermediate variable I by means of upper/lower triangular function solver, and outputting the intermediate variable I to the block code encoding unit (208);
the block code encoding unit (208), being used for according to $C=I \times G_{ldgc}$, using the generator matrix G_{ldgc} to encode the intermediate variable I, and generating an encoding sequence with a length of N+L-K; and
a bit deleting unit (210), being used for deleting L-K known bits from the encoding sequence with a length of N+L-K, generating an encoding sequence with a length of N.

9. The encoding device according to claim 8, wherein, the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} being an upper-left triangular matrix, or a lower-left triangular matrix, or an upper-right triangular matrix, or a lower-right triangular matrix.

10. The encoding device according to claim 9, wherein, when the phalanx $G_{ldgc}(1:L, 1:L)$ that consists of L rows and the first L columns of the generator matrix G_{ldgc} is an upper-left triangular matrix or a lower-left triangular matrix, the bit filling unit adding L-K known bits in front of the K-length information bit sequence that is to be encoded.

11. The encoding device according to claim 9, wherein, when the phalanx $G_{ldgc}(1:L, 1:L)$ that consists L rows and the first L columns of the generator matrix G_{ldgc} is an upper-right triangular matrix or a lower-right triangular matrix, the bit filling unit adding L-K known bits after the K-length information bit sequence that is to be encoded.

12. The encoding device according to any of claims 8 to 11, wherein, a column weight of the generator matrix G_{ldgc} being in accordance with a degree distribution principle similar to that of LT codes.

13. A device for encoding the Low Density Generator Matrix Code, LDGC, comprising:

a matrix generating unit (602), being used for generating a generator matrix G_{ldgc} with K rows and N columns, wherein, a phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} being an upper triangular matrix or a lower triangular matrix, the K and N being positive integers, and $K<N$;
a precoding unit (604), being used for according to $I \times G_{ldgc}(1:K, 1:K)=s$, by using a K-length information bit sequence s that is to be encoded and the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc}, generating an intermediate variable I by means of upper/lower triangular function solver; and
a block code encoding unit (606), being used for according to $C=I \times G_{ldgc}$, using the generator matrix G_{ldgc} to encode the intermediate variable I, and generating a code word with a length of N.

14. The encoding device according to claim 13, wherein, the phalanx $G_{ldgc}(1:K, 1:K)$ that consists of K rows and the first K columns of the generator matrix G_{ldgc} being an upper-left triangular matrix, or a lower-left triangular matrix, or an upper-right triangular matrix, or a lower-right triangular matrix.

15. The encoding device according to claim 13 or claim 14, wherein, a column weight of the generator matrix G_{ldgc} being

in accordance with a degree distribution principle similar to that of LT codes.

Patentansprüche

1. Verfahren zum Codieren des Low Density Generator Matrix Codes, LDGC, mit den folgenden Schritten:

   Bilden (S102) einer Generatormatrix G_{ldgc} mit L Zeilen und N+L-K Spalten, wobei eine Phalanx G_{ldgc}(1:L, 1 :L), die aus L Zeilen und die ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere Dreiecksmatrix oder eine untere Dreiecksmatrix ist, wobei K, L und N positive ganze Zahlen sind und K <L <N ist;
   Hinzufügen (S104) L-K bekannter Bits zu einer K-Längen Informationsbitsequenz, die codiert werden soll, und Erzeugen einer Informationsbitfolge m mit einer Länge von L;
   gemäß $I*G_{ldgc}(1:L, 1 :L) = m$, wobei die Phalanx G_{ldgc}(1:L, 1 :L) aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} und der Informationsbitfolge m mit einer Länge von L verwendet wird, wobei eine Zwischenvariable I mittels eines oberen/unteren Dreiecksfunktionslösers erzeugt wird und nach $C=I*G_{ldgc}$, unter Verwendung der Generatormatrix G_{ldgc}, um die Zwischenvariable I zu codieren, und eine Codierungssequenz mit der Länge von N+L-K (S106) erzeugt wird; und
   Löschen von L-K bekannten Bits aus der Codierungssequenz mit einer Länge von N+L-K und Erzeugen einer Codierungssequenz mit einer Länge von N (S108).

2. Kodierverfahren nach Anspruch 1, wobei, wenn die Phalanx G_{ldgc} (1 :L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere linke Dreiecksmatrix oder eine untere linke Dreiecksmatrix ist, die L-K bekannten Bits vor der zu codierenden K-Längen Informationsbitfolge hinzugefügt werden.

3. Kodierverfahren nach Anspruch 1, wobei, wenn die Phalanx G_{ldgc} (1 :L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere rechte Dreiecksmatrix oder eine untere rechte Dreiecksmatrix ist, die L-K bekannten Bits nach der zu codierenden K-Längen Informationsbitfolge hinzugefügt werden.

4. Kodierverfahren nach einem der Ansprüche 1 bis 3, wobei ein Spaltengewicht der Generatormatrix G_{ldgc} in Übereinstimmung mit einem Gradverteilungsprinzip ähnlich dem von LT Codes ist.

5. Verfahren zum Codieren des Low Density Generator Matrix Codes, LDGC, mit den folgenden Schritten:

   Bilden (S502) einer Generatormatrix G_{ldgc} mit K Zeilen und N Spalten, wobei eine Phalanx G_{ldgc}(1:K,1:K), die aus K Zeilen und den ersten K Spalten der Generatormatrix G_{ldgc} besteht, eine obere Dreiecksmatrix oder eine untere Dreiecksmatrix ist, wobei K und N positive ganze Zahlen sind und K < N ist;
   gemäß $I*G_{ldgc}(1:K, 1 :K)=s$ unter Verwendung der zu codierenden K-Längen Informationsbitfolge s und der Phalanx G_{ldgc}(1:K; 1:K), bestehend aus K Reihen und den ersten K Spalten der Generatormatrix G_{ldgc}, Erzeugen (S504) einer Zwischenvariablen I mittels eines oberen/unteren Dreiecksfunktionslösers wird; und
   gemäß $C=I*G_{ldgc}$, die Zwischenvariable I unter Verwendung der Generatormatrix G_{ldgc} codiert und ein Codewort mit einer Länge von N (S506) erzeugt wird.

6. Kodierverfahren nach Anspruch 5, wobei die Phalanx G_{ldgc} (1 :K, 1 :K), die aus K Zeilen und den ersten K Spalten der Generatormatrix G_{ldgc} besteht, eine obere linke Dreiecksmatrix, oder eine untere linke Dreiecksmatrix, oder eine obere rechte Dreiecksmatrix, oder eine untere rechte Dreiecksmatrix ist.

7. Kodierverfahren nach Anspruch 5 oder 6, wobei ein Spaltengewicht der Generatormatrix G_{ldgc} in Übereinstimmung mit einem Gradverteilungsprinzip ähnlich dem von LT Codes ist.

8. Vorrichtung für die Codierung des Low Density Generator Matrix Codes, LDGC, umfassend:

   eine Matrixerzeugungseinheit (202), die zur Erzeugung einer Generatormatrix G_{ldgc} mit L Zeilen und N+L-K Spalten verwendet wird, die die erzeugte Matrix G_{ldgc} an eine Blockcode-Codiereinheit (208) ausgibt und eine Phalanx G_{ldgc}(1:L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, an eine Vorcodiereinheit (206) ausgibt, wobei die Phalanx G_{ldgc}(1:L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere Dreiecksmatrix oder eine untere Dreiecksmatrix ist, wobei K, L und N positive ganze Zahlen sind und K <L <N ist;
   eine Bit-Fülleinheit (204), die zum Hinzufügen von L-K bekannten Bits zu einer K-Längen Informationsbitfolge,

die codiert werden soll, verwendet wird, erzeugt eine Informationsbitfolge mit einer Länge von L , und Ausgeben der erzeugten Informationsbitfolge m mit einer Länge von L zu der Vorcodiereinheit (206);
wobei die Vorcodiereinheit (206), welche nach $I^* G_{ldgc}(1:L, 1:L)=m$ verwendet wird, indem die Phalanx $G_{ldgc}(1:L, 1:L)$ verwendet wird, die aus L Reihen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht und der Informationsbitfolge m mit einer Länge von L, eine Zwischenvariable I mittels eines oberen/unteren Dreiecksfunktionslösers erzeugt und die Zwischenvariable I an die Blockcode-Codiereinheit (208) ausgibt;
wobei die Blockcode-Codiereinheit (208) gemäß $C=I^* G_{ldgc}$ verwendet wird, wobei die Generatormatrix G_{ldgc} verwendet wird, um die Zwischenvariable I zu codieren und eine Codierungssequenz mit einer Länge von N+L-K zu erzeugen; und
eine Bit-Löscheinheit (210), die zum Löschen von L-K bekannten Bits aus der Codierungssequenz mit einer Länge von N+L-K verwendet wird, wodurch eine Codierungssequenz mit einer Länge von N erzeugt wird.

9. Kodiervorrichtung nach Anspruch 8, wobei die Phalanx G_{ldgc} (1 :L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere linke Dreiecksmatrix, oder eine untere linke Dreiecksmatrix, oder eine obere rechte Dreiecksmatrix, oder eine untere rechte Dreiecksmatrix ist.

10. Kodiervorrichtung nach Anspruch 9, wobei, wenn die Phalanx G_{ldgc} (1 :L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere linke Dreiecksmatrix oder eine untere linke Dreiecksmatrix ist, die Bit-Fülleinheit die L-K bekannten Bits vor der zu codierenden K-Längen Informationsbitfolge hinzufügt.

11. Kodiervorrichtung nach Anspruch 9, wobei, wenn die Phalanx G_{ldgc} (1 :L, 1 :L), die aus L Zeilen und den ersten L Spalten der Generatormatrix G_{ldgc} besteht, eine obere rechte Dreiecksmatrix oder eine untere rechte Dreiecksmatrix ist, die Bit-Fülleinheit die L-K bekannten Bits nach der zu codierenden K-Längen Informationsbitfolge hinzufügt.

12. Kodiervorrichtung nach einem der Ansprüche 8 bis 11, wobei ein Spaltengewicht der Generatormatrix G_{ldgc} in Übereinstimmung mit einem Gradverteilungsprinzip ähnlich dem von LT Codes ist.

13. Vorrichtung zum Codieren des Low Density Generator Matrix Codes, LDGC, umfassend:

eine Matrixerzeugungseinheit (602), die zur Erzeugung einer Generatormatrix G_{ldgc} mit K Zeilen und N Spalten verwendet wird, wobei eine Phalanx $G_{ldgc}(1:K, 1:K)$, die aus K Zeilen und den ersten K Spalten der Generatormatrix besteht, wobei G_{ldgc} eine obere Dreiecksmatrix oder eine untere Dreiecksmatrix ist, wobei K und N positive ganze Zahlen sind und K <N ist;
eine Vorcodiereinheit (604), die nach $I^*G_{ldgc}(1:K, 1:K)=s$ verwendet wird, indem eine zu codierende K-Längen Informationsbitfolge s verwendet wird und die Phalanx $G_{ldgc}(1:K, 1:K)$, die aus K Zeilen und den ersten K Spalten der Generatormatrix G_{ldgc} besteht, die eine Zwischenvariable I mit Hilfe eines oberen/unteren Dreiecksfunktionslösers erzeugen; und
eine Blockcode-Codiereinheit (606), die gemäß $C=I^*G_{ldgc}$ verwendet wird, wobei die Generatormatrix G_{ldgc} verwendet wird, um die Zwischenvariable I zu codieren und ein Codewort mit einer Länge von N zu erzeugen.

14. Vorrichtung zum Codieren nach Anspruch 13, wobei die Phalanx G_{ldgc} (1 :K, 1 :K), die aus K Zeilen und den ersten K Spalten der Generatormatrix G_{ldgc} besteht, eine obere linke Dreiecksmatrix, oder eine untere linke Dreiecksmatrix, oder eine obere rechte Dreiecksmatrix, oder eine untere rechte Dreiecksmatrix ist.

15. Vorrichtung zum Codieren nach Anspruch 13 oder 14, wobei ein Spaltengewicht der Generatormatrix G_{ldgc} in Übereinstimmung mit einem Gradverteilungsprinzip ähnlich dem von LT Codes ist.

Revendications

1. Procédé pour encoder le code de matrice génératrice de faible densité, LDGC, comprenant les étapes suivantes :

construire (S102) une matrice génératrice G_{ldgc} avec L rangées et N+L-K colonnes, dans lequel une phalange $G_{ldgc}(1:L, 1:L)$ composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure ou une matrice triangulaire inférieure, K, L et N étant des entiers positifs, et K < L < N ;
ajouter (S104) L-K bits connus à une séquence de bits d'informations de longueur K qui doit être encodée, et générer une séquence de bits d'informations m avec une longueur L ;

conformément à I x G_{ldgc}(1:L, 1 :L) = m, en utilisant la phalange G_{ldgc}(1:L, 1:L) composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc}, et la séquence de bits d'informations m avec une longueur L, générer une variable I intermédiaire au moyen d'un résolveur de fonction triangulaire supérieure/inférieure, et conformément à C = I x G_{ldgc}, en utilisant la matrice génératrice G_{ldgc} pour encoder la variable I intermédiaire, et générer une séquence d'encodage avec une longueur N+L-K (S106) ; et
effacer les L-K bits connus de la séquence d'encodage avec une longueur N+L-K, et générer une séquence d'encodage avec une longueur N (S108).

2. Procédé d'encodage selon la revendication 1, comprenant, lorsque la phalange G_{ldgc}(1:L, 1 :L) qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure gauche ou une matrice triangulaire inférieure gauche, ajouter les L-K bits connus à l'avant de la séquence de bits d'informations de longueur K qui doit être encodée.

3. Procédé d'encodage selon la revendication 1, dans lequel, lorsque la phalange G_{ldgc}(1:L, 1 :L) qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure droite ou une matrice triangulaire inférieure droite, ajouter les L-K bits connus après la séquence de bits d'informations de longueur K qui doit être encodée.

4. Procédé d'encodage selon l'une quelconque des revendications 1 à 3, dans lequel un poids de colonne de la matrice génératrice G_{ldgc} est conforme à un principe de distribution de degrés similaire à celui des codes LT.

5. Procédé pour l'encodage du code de matrice génératrice de faible densité, LDGC, comprenant les étapes suivantes :

construire (S502) une matrice génératrice G_{ldgc} avec K rangées et N colonnes, dans lequel une phalange G_{ldgc}(1:K, 1 :K) qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure ou une matrice triangulaire inférieure, K et N étant des entiers positifs, et K < N ;
conformément à I x G_{ldgc}(1:K, 1 :K) = s, en utilisant la séquence de bits d'informations de longueur K s qui doit être encodée et la phalange G_{ldgc}(1:K, 1 :K) qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc}, générer (S504) une variable I intermédiaire au moyen d'un résolveur de fonction triangulaire supérieure/inférieure ; et
conformément à C = I x G_{ldgc}, encoder la variable I intermédiaire en utilisant la matrice génératrice G_{ldgc}, et générer un mot de code avec une longueur N (S506).

6. Procédé d'encodage selon la revendication 5, dans lequel la phalange G_{ldgc}(1:K, 1 :K) qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure gauche, ou une matrice triangulaire inférieure gauche, ou une matrice triangulaire supérieure droite, ou une matrice triangulaire inférieure droite.

7. Procédé d'encodage selon la revendication 5 ou la revendication 6, dans lequel un poids de colonne de la matrice génératrice G_{ldgc} est conforme à un principe de distribution de degrés similaire à celui des codes LT.

8. Dispositif pour l'encodage du code de matrice génératrice de faible densité, G_{ldgc}, comprenant :

une unité de génération de matrice (202), utilisée pour générer une matrice génératrice G_{ldgc} avec L rangées et N+L-K colonnes, sortir la matrice G_{ldgc} générée vers une unité d'encodage de code de bloc (208), et sortir une phalange G_{ldgc}(1:L, 1 :L) qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} vers une unité de précodage (206), dans lequel, la phalange G_{ldgc}(1:L, 1 :L) qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure ou une matrice triangulaire inférieure, K, L et N étant des entiers positifs, et K < L < N ;
une unité de remplissage avec des bits (204), utilisée pour ajouter L-K bits connus à une séquence de bits d'informations de longueur K qui doit être encodée, générer une séquence de bits d'informations m avec une longueur L, et sortir la séquence de bits d'informations m générée avec une longueur L vers l'unité de précodage (206) ;
l'unité de précodage (206), utilisée, conformément à I x G_{ldgc}(1:L, 1 :L) = m, en utilisant la phalange G_{ldgc}(1:L, 1 :L) qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} et la séquence de bits d'informations m avec une longueur L, générer une variable I intermédiaire au moyen d'un résolveur de fonction triangulaire supérieure/inférieure, et sortir la variable I intermédiaire vers l'unité d'encodage de code

de bloc (208) ;
l'unité d'encodage de code de bloc (208), utilisée, conformément à $C = I \times G_{ldgc}$, en utilisant la matrice génératrice G_{ldgc} pour encoder la variable I intermédiaire, et générer une séquence d'encodage avec une longueur N+L-K ; et une unité d'effacement de bit (210), utilisée pour effacer L-K bits connus de la séquence d'encodage avec une longueur N+L-K, et générer une séquence d'encodage avec une longueur N.

9. Dispositif d'encodage selon la revendication 8, dans lequel la phalange $G_{ldgc}(1:L, 1:L)$ qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure gauche, ou une matrice triangulaire inférieure gauche, ou une matrice triangulaire supérieure droite, ou une matrice triangulaire inférieure droite.

10. Dispositif d'encodage selon la revendication 9, dans lequel, lorsque la phalange $G_{ldgc}(1:L, 1:L)$ qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure gauche ou une matrice triangulaire inférieure gauche, l'unité de remplissage avec des bits ajoute L-K bits connus à l'avant de la séquence de bits d'informations de longueur K qui doit être encodée.

11. Dispositif d'encodage selon la revendication 9, dans lequel, lorsque la phalange $G_{ldgc}(1:L, 1:L)$ qui est composée de L rangées et des L premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure droite ou une matrice triangulaire inférieure droite, l'unité de remplissage avec des bits ajoute L-K bits connus après la séquence de bits d'informations de longueur K qui doit être encodée.

12. Dispositif d'encodage selon l'une quelconque des revendications 8 à 11, dans lequel un poids de colonne de la matrice génératrice G_{ldgc} est conforme à un principe de distribution de degrés similaire à celui des codes LT.

13. Dispositif pour encoder le code de matrice génératrice de faible densité, LDGC, comprenant :

une unité de génération de matrice (602), utilisée pour générer une matrice génératrice G_{ldgc} avec K rangées et N colonnes, dans lequel une phalange $G_{ldgc}(1:K, 1:K)$ qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure ou une matrice triangulaire inférieure, K et N étant des entiers positifs, et $K < N$;
une unité de précodage (604), utilisée, conformément à $I \times G_{ldgc}(1:K, 1:K) = s$, en utilisant une séquence de bits d'informations de longueur K qui doit être encodée et la phalange $G_{ldgc}(1:K, 1:K)$ qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc}, générer une variable I intermédiaire au moyen d'un résolveur de fonction triangulaire supérieure/inférieure ; et
une unité d'encodage de code de bloc (606), utilisée, conformément à $C = I \times G_{ldgc}$, en utilisant la matrice génératrice G_{ldgc} pour encoder la variable I intermédiaire, et générer un mot de code avec une longueur N.

14. Dispositif d'encodage selon la revendication 13, dans lequel la phalange $G_{ldgc}(1:K, 1:K)$ qui est composée de K rangées et des K premières colonnes de la matrice génératrice G_{ldgc} est une matrice triangulaire supérieure gauche, ou une matrice triangulaire inférieure gauche, ou une matrice triangulaire supérieure droite, ou une matrice triangulaire inférieure droite.

15. Dispositif d'encodage selon la revendication 13 ou la revendication 14, dans lequel un poids de colonne de la matrice génératrice G_{ldgc} est conforme à un principe de distribution de degrés similaire à celui des codes LT.

Fig. 1

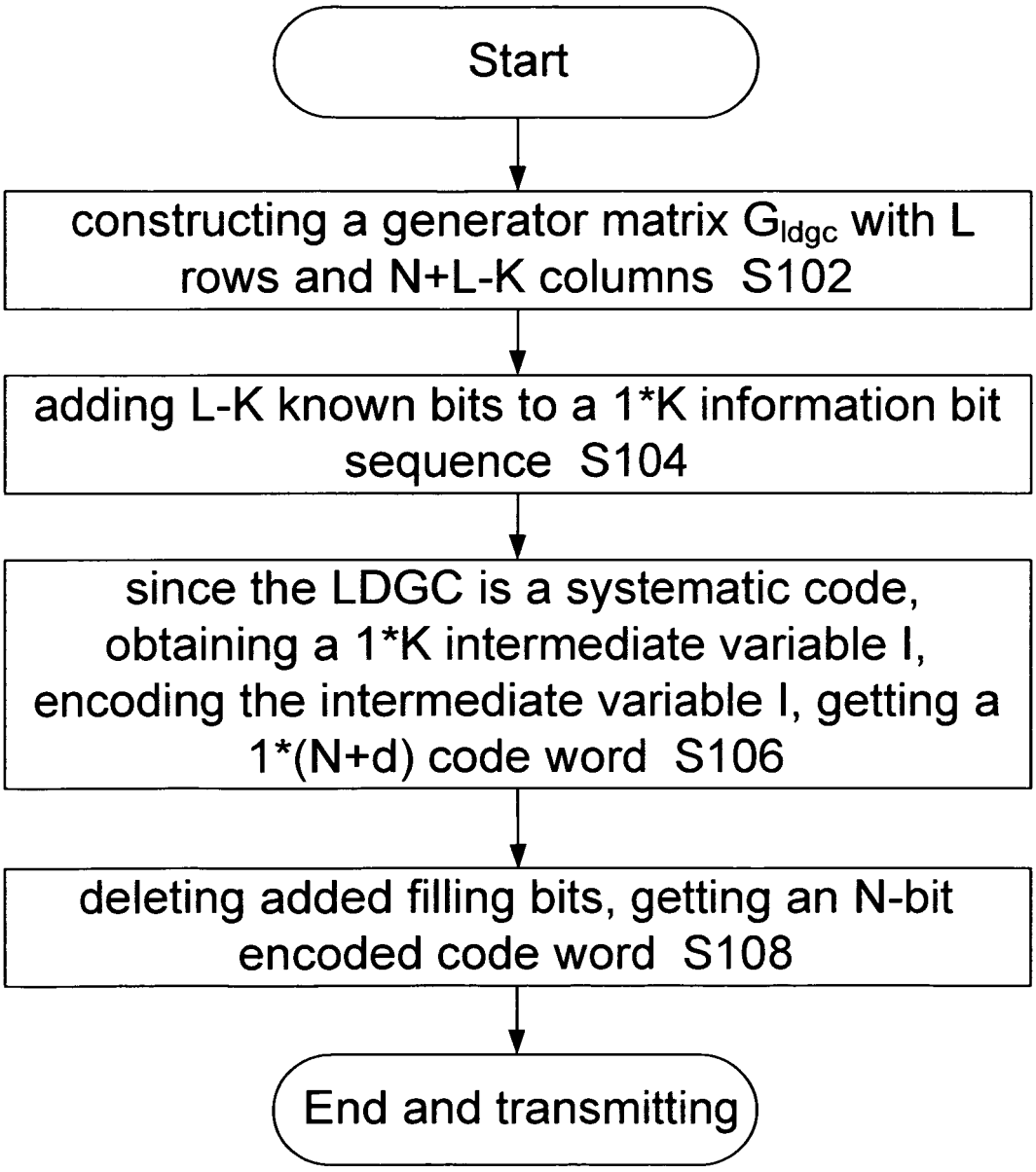

Start
constructing a generator matrix G_{ldgc} with L rows and N+L-K columns S102
adding L-K known bits to a 1*K information bit sequence S104
since the LDGC is a systematic code, obtaining a 1*K intermediate variable I, encoding the intermediate variable I, getting a 1*(N+d) code word S106
deleting added filling bits, getting an N-bit encoded code word S108
End and transmitting

Fig. 2

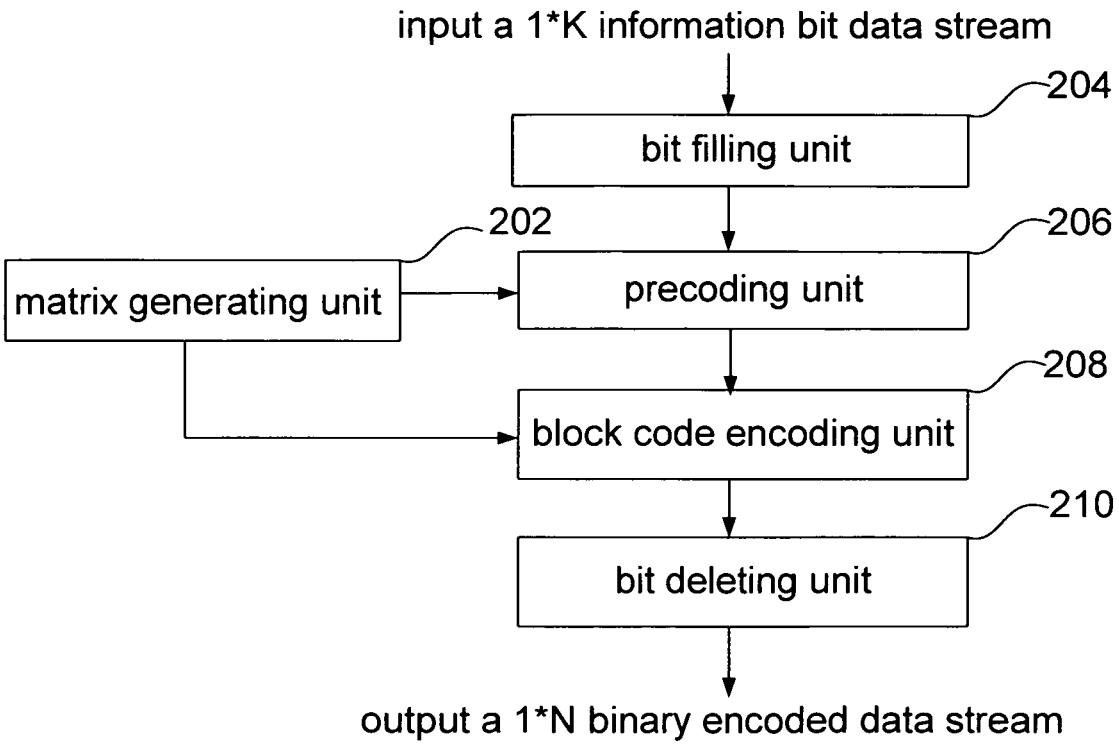

input a 1*K information bit data stream
204
bit filling unit
202
206
matrix generating unit
precoding unit
208
block code encoding unit
210
bit deleting unit
output a 1*N binary encoded data stream

Fig. 3

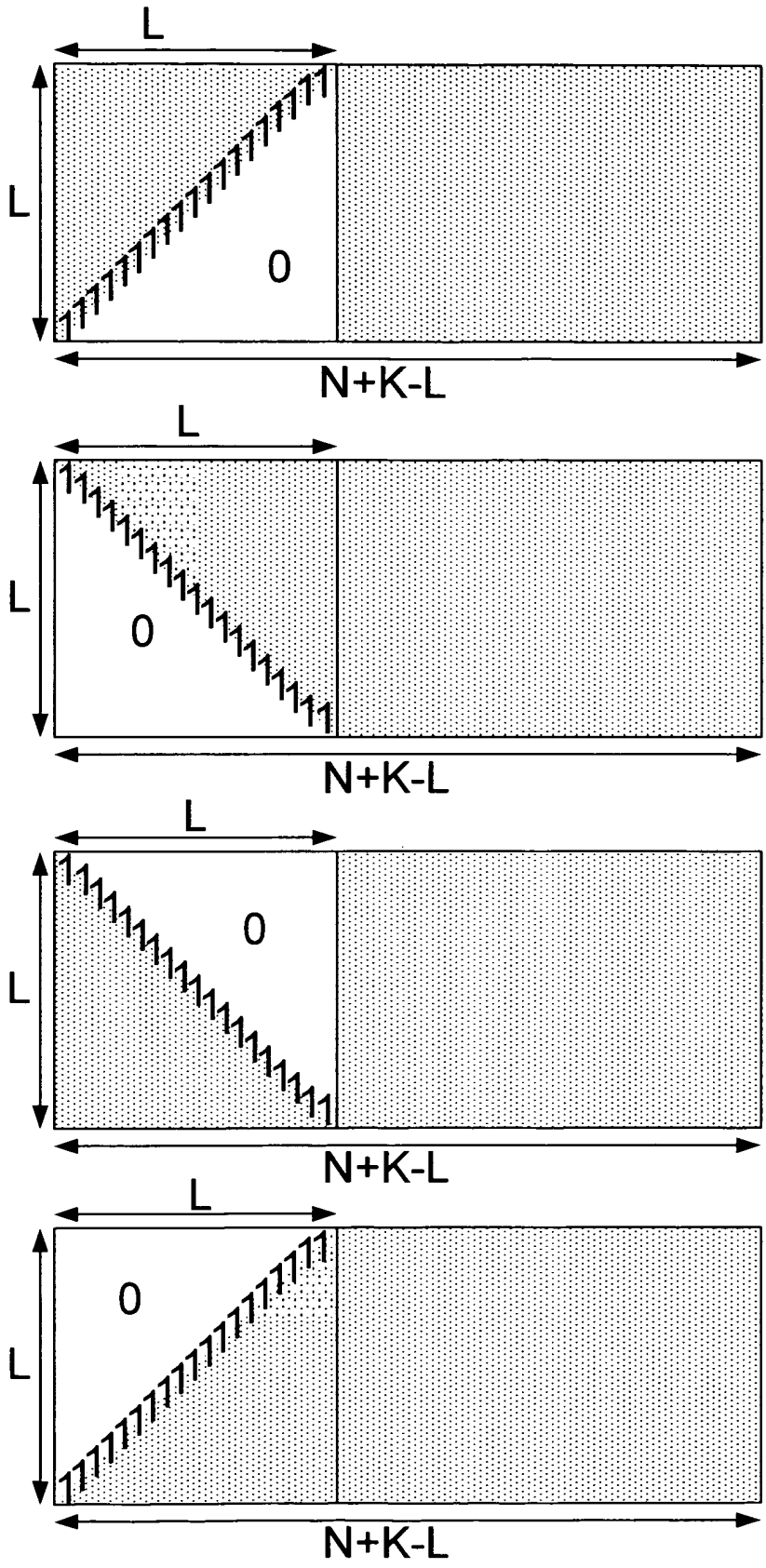


Fig. 4

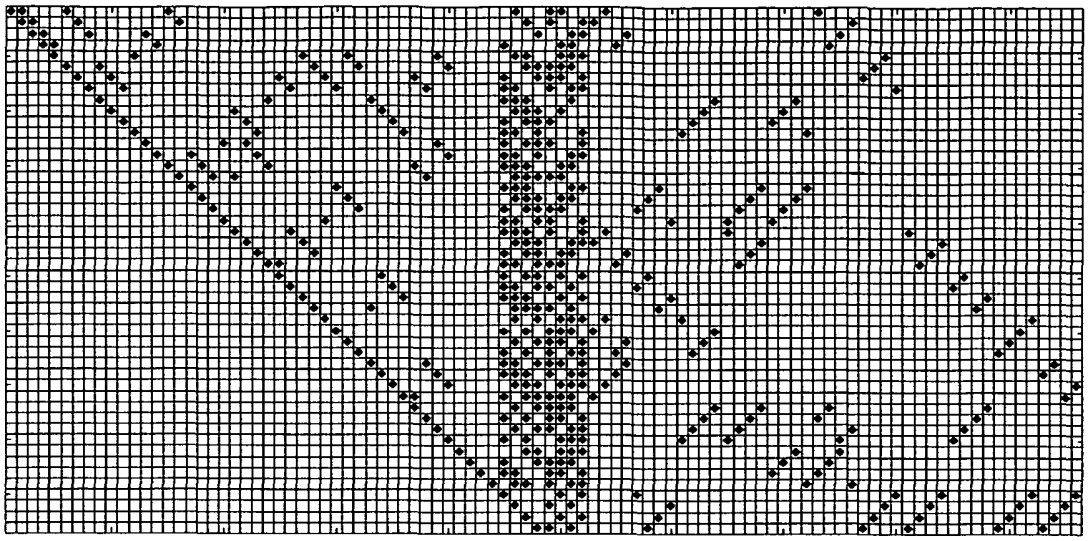


Fig. 5

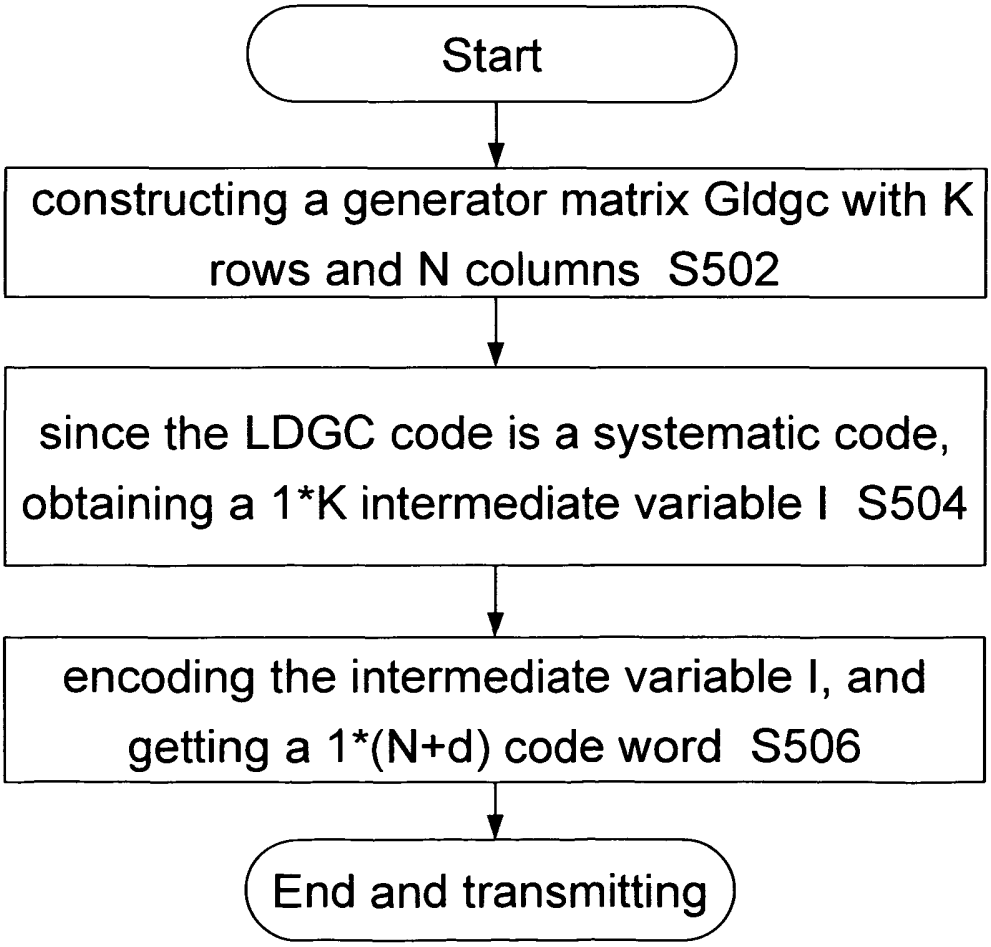

Start
constructing a generator matrix Gldgc with K rows and N columns S502
since the LDGC code is a systematic code, obtaining a 1*K intermediate variable I S504
encoding the intermediate variable I, and getting a 1*(N+d) code word S506
End and transmitting

Fig. 6

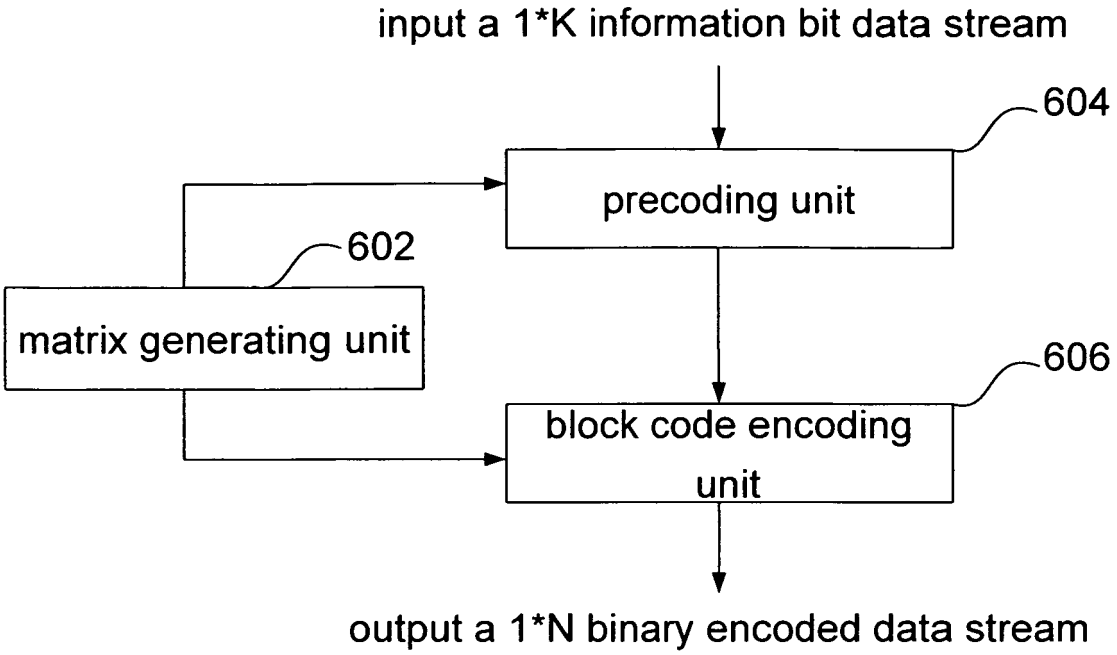

input a 1*K information bit data stream
604
precoding unit
602
matrix generating unit
606
block code encoding unit
output a 1*N binary encoded data stream

Fig. 7

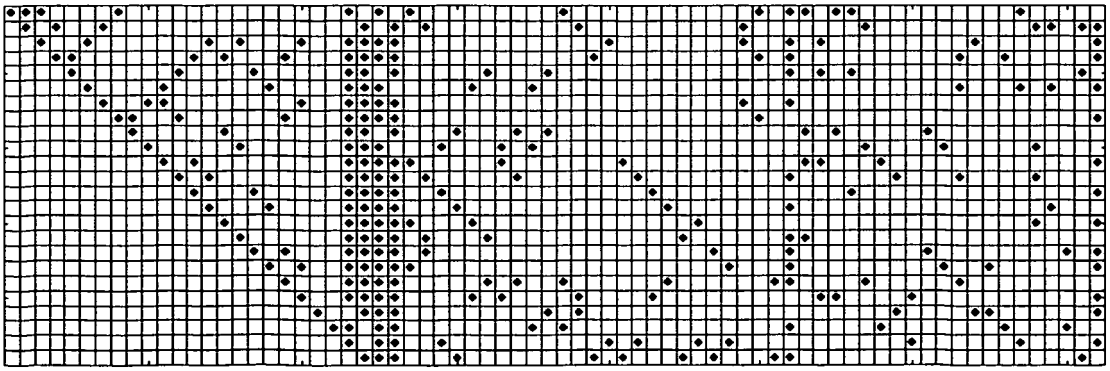
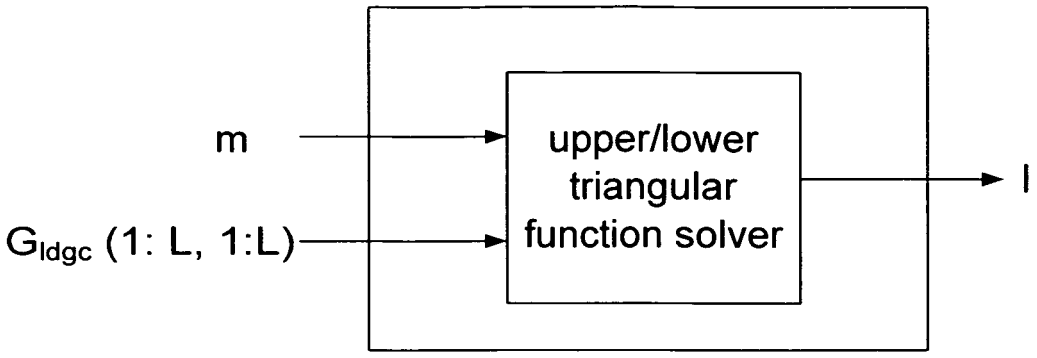
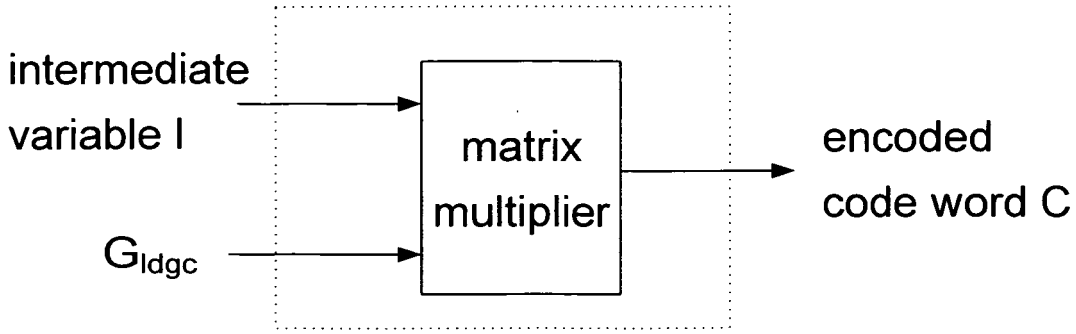


Fig. 8



precoding unit

Fig. 9



block code encoding unit

REFERENCES CITED IN THE DESCRIPTION

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Non-patent literature cited in the description

- **VINCENT ROCA ; CHRISTOPH NEUMANN.** *Design, Evaluation and Comparison of Four Large Block FEC Codes, LDPC, LDGM, LDGM Staircase and LDGM Triangle,* 09 June 2004, http://hal **[0009]**