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(54) **CONTROL DEVICE AND SYSTEM PROGRAM, AND RECORDING MEDIUM**

STEUERVORRICHTUNG UND SYSTEMPROGRAMM SOWIE AUFZEICHNUNGSMEDIUM

DISPOSITIF DE COMMANDE ET PROGRAMME-SYSTÈME, ET SUPPORT D'ENREGISTREMENT

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(73) Proprietor: **Omron Corporation**
Kyoto-shi, Kyoto 600-8530 (JP)

(72) Inventors:
• **HAMASAKI, Osamu**
Kyoto-shi
Kyoto 600-8530 (JP)
• **EGUCHI, Shigeyuki**
Kyoto-shi
Kyoto 600-8530 (JP)
• **KOBAYASHI, Akiro**
Kyoto-shi
Kyoto 600-8530 (JP)

• **INAME, Yukio**
Kyoto-shi
Kyoto 600-8530 (JP)
• **YAOITA, Koji**
Kyoto-shi
Kyoto 600-8530 (JP)

(74) Representative: **Horn Kleimann Waitzhofer**
Patentanwälte PartG mbB
Ganghoferstrasse 29a
80339 München (DE)

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Description**TECHNICAL FIELD**

[0001] The present invention relates to a control device for controlling a control object apparatus, a system program executed by the control device and a recording medium in which the system program is recorded.

BACKGROUND ART

[0002] Conventionally, a PLC (Programmable Logic Controller) for controlling an industrial machine or the like installed in a production line of a factory has been known.

[0003] The PLC is provided with a CPU unit for controlling the PLC, an input unit to which a signal from a sensor is inputted, and an output unit for outputting a control signal to an industrial machine or the like. The CPU unit has a memory for storing a user program for controlling the industrial machine or the like. This user program is configured so as to be editable by a tool device connected to the CPU unit.

[0004] Moreover, in the PLC, a process for acquiring a signal inputted to the input unit in the memory of the CPU unit, a process for executing the user program, a process for writing execution results (operation results) of the user program in the memory and for sending them to the output unit, and a peripheral process for transmitting and receiving data to and from the tool device are executed repeatedly. Thus, the PLC can control an industrial machine or the like based upon an input from the sensor.

[0005] Moreover, conventionally, a PAC (Programmable Automation Controller) provided with a high degree software function of a personal computer and reliability of the PLC has been known (for example, see Patent Document 1 (Japanese Unexamined Patent Publication No. 2009-181443)).

[0006] The PAC of the above-mentioned Patent Document 1 (Japanese Unexamined Patent Publication No. 2009-181443) is configured to execute a plurality of tasks in parallel with one another in a time sharing manner, with each task being executed in cycles based on each task. That is, this PAC carries out multitask processes.

[0007] In this PAC, it is proposed that an operation mode for executing a user program in each task and a program mode for executing no user program in each task are prepared in a switchable manner.

[0008] In this case, the operation mode is a mode for actually controlling an industrial machine or the like by executing a user program, and the program mode is a mode which allows the user to edit a user program by using a tool device, without executing a user program. For this reason, the execution of a user program is stopped by switching modes from the operation mode to the program mode, while the execution of the user program is started by switching modes from the program mode to the operation mode.

PRIOR ART DOCUMENT**PATENT DOCUMENT**

[0009] Patent Document 1: Japanese Unexamined Patent Publication No. 2009-181443

[0010] Jakob Engblom, "Debugging Real-Time Multiprocessor Systems", Class #249/269, Embedded Systems Conference, Silicon Valley 2007, discusses ways to debug parallel software running on top of multiprocessor hardware, and the types of errors that occur due to parallelism.

SUMMARY OF THE INVENTION**PROBLEMS TO BE SOLVED BY THE INVENTION**

[0011] In the above-mentioned PAC, however, upon carrying out multitask processes, a problem arises in that it becomes difficult to appropriately carry out switching of modes because cycles of the respective tasks are different from one another. More specifically, even though an execution of a user program of a task having a short cycle is completed in switching modes from an operation mode to a program mode, when the switching of mode is carried out during an execution of a user program of a task having a long cycle, a problem arises in that a normal operation result is not obtained in the user program having a long cycle, with the result that the PAC fails to accurately control an industrial machine or the like.

[0012] The present invention has been devised to solve the above-mentioned problem, and the object of the present invention is to provide a control device that can appropriately switch operating modes even when multitask processes are being carried out and a system program executed in the control device, as well as a recording medium in which the system program is recorded.

MEANS FOR SOLVING THE PROBLEM

[0013] A control device according to one aspect of the present invention is a control device for controlling a control object apparatus including: a task executing unit configured to execute a plurality of tasks in parallel with one another, and to execute each task in a cycle corresponding to the task; and a mode switching unit configured to switch operating modes of the task executing unit. The task executing unit is configured such that in a first mode, a user program is executed for each task, and in a second mode, a user program is not executed for each task. The mode switching unit is configured such that the operating mode of the task executing unit is switched when the plurality of tasks, which are being executed by the task executing unit, are synchronized with one another.

[0014] Additionally, to execute a plurality of tasks in parallel with one another includes a case in which a plurality of tasks are executed simultaneously in parallel with

one another and a case in which a plurality of tasks are executed in parallel with one another in a time sharing manner. Moreover, the time at which tasks are synchronized with one another refers to a point of time after executions of a plurality of tasks carried out in cycles based on each task and prior to the simultaneous starts of the execution cycles of the plurality of tasks. Furthermore, the user program is a program for controlling a control object apparatus, and a first mode is an operating mode in which the control on the control object apparatus is carried out, and a second mode is an operating mode in which no control on the control object apparatus is carried out.

[0015] With this configuration, in the case that a plurality of tasks, which are being executed, are synchronized with one another, switching operating modes in the task executing unit can prevent the execution of a user program for each task from being interrupted even if the cycles of the respective tasks are different from one another. Therefore, it is possible to appropriately switch operating modes when multitask processes are carried out. As a result, since a correct operation result can be obtained in the user program for each task, the control device is allowed to normally control the control object apparatus.

[0016] The control device may include a management unit configured to manage executions of a plurality of tasks by the task executing unit, and to manage timing in which the mode switching unit switches the operating mode in the task executing unit.

[0017] With this configuration, operating modes in the task executing unit can be easily switched by the management unit, when the plurality of tasks, which are being executed, are synchronized with one another.

[0018] The control device including the management unit includes a receiving unit configured to receive a switching instruction for operating modes in the task executing unit. The management unit may be configured such that when the plurality of tasks, which are being executed by the task executing unit, are synchronized with one another after the receipt of the switching instruction for operating modes by the receiving unit, the operating mode in the task executing unit is switched by the mode switching unit.

[0019] With this configuration, the operating modes in the task executing unit can be switched, when the plurality of tasks, which are being executed, are synchronized with one another after a receipt of a switching instruction.

[0020] In this case, the control device includes a storing unit configured to switch a user program, and a communication unit configured to communicate with a tool device for editing the user program stored in the storing unit. The receiving unit may be configured to receive the switching instruction transmitted from the tool device via the communication unit.

[0021] With this configuration, the user is allowed to switch operating modes in the task execution unit by us-

ing the tool device.

[0022] In the control device, the task executing unit may be configured so as to execute a plurality of tasks simultaneously in parallel with one another.

[0023] With this configuration, it is possible to improve the task processing capability.

[0024] In the control device, the task executing unit may be configured so as to execute a plurality of tasks in parallel with one another in a time sharing manner.

[0025] With this configuration, it is possible to carry out the task processing efficiently.

[0026] In the control device, a plurality of tasks may include a reference task and tasks other than the reference task, with each of the tasks other than the reference task being set to have a cycle having an integral multiple of the cycle of the reference task.

[0027] With this configuration, it is possible to easily make a plurality of tasks synchronized with one another.

[0028] A system program according to another aspect of the present invention is a system program, which is executed by a control device including a first mode in which a control object apparatus is controlled by executing a user program in each task and a second mode in which no user program is executed in each task, the system program including: a sequence for executing a plurality of tasks in parallel in a cycle based on each task; and a sequence for switching modes between the first mode and the second mode when the plural tasks, which are being executed, are synchronized with one another.

[0029] With this configuration, it is possible to appropriately switch operating modes while multitask processes are being carried out by the control device.

[0030] A recording media of another aspect of the present invention is a recording media that has recorded the system program, allowing a computer to read the system program.

EFFECT OF THE INVENTION

[0031] The present invention can provide a control device capable of appropriately switching operating modes when multitask processes are being performed, a system program to be executed in such a control device, and a recording medium in which the system program is recorded.

BRIEF DESCRIPTION OF THE DRAWINGS

[0032]

Fig. 1 is a hardware block diagram showing a configuration of a PAC provided with a CPU unit in accordance with a first embodiment of the present invention.

Fig. 2 is a hardware block diagram showing configurations of the CPU unit of the PAC shown in Fig. 1. Fig. 3 is a functional block diagram of the CPU unit shown in Fig. 2.

Fig. 4 is a view for explaining information relating to tasks stored in a ROM of the CPU unit shown in Fig. 2.

Fig. 5 shows a time chart for explaining operations, at the time of an operation mode, of the CPU unit in accordance with the first embodiment of the present invention.

Fig. 6 shows a time chart for explaining operations, at the time of a program mode, of the CPU unit in accordance with the first embodiment of the present invention.

Fig. 7 shows a time chart for explaining operations, upon switching modes from the operation mode to the program mode, of the CPU unit in accordance with the first embodiment of the present invention.

Fig. 8 shows a time chart for explaining operations, upon switching modes from the program mode to the operation mode, of the CPU unit in accordance with the first embodiment of the present invention.

Fig. 9 is a hardware block diagram showing configurations of a CPU unit in accordance with a second embodiment of the present invention.

Fig. 10 is a view for explaining information relating to tasks stored in a ROM of the CPU unit shown in Fig. 9.

Fig. 11 shows a time chart for explaining operations, at the time of an operation mode, of the CPU unit in accordance with the second embodiment of the present invention.

MODE FOR CARRYING OUT THE INVENTION

[0033] Hereinafter, referring to drawings, the following description will be given of an embodiment of the present invention.

<First Embodiment>

[0034] First, referring to Figs. 1 to 4, the following description will be given of a configuration of a PAC 100 provided with a CPU unit 1 in accordance with a first embodiment.

[0035] As shown in Fig. 1, the PAC 100 is provided with a CPU unit 1 for controlling the PAC 100, an input unit 2 to which a signal from a sensor 50 is inputted, an output unit 3 for outputting a control signal to a control object apparatus 60 such as an industrial machine, a communication unit 4 for communicating with another PAC 70, a power supply unit 5 for supplying a power to the respective units and an internal bus 6 for connecting the respective units to one another. Additionally, the CPU unit 1 is one example of "the control device" of the present invention.

[0036] The PAC 100 is configured to control the control object apparatus 60 based upon an input from the sensor 50 by executing a task. In this case, the task includes the execution of a user program for controlling the control object apparatus 60.

[0037] The CPU unit 1 in which a system program 121 (see Fig. 2) is installed is configured so that operations of the CPU unit 1 are controlled by executing the system program 121. In this case, the system program 121 includes a program for executing an acquiring process of a signal inputted to the input unit 2 into the CPU unit 1, a program for executing a sending process of operation results by the CPU unit 1 to the output unit 3, and the like.

[0038] As shown in Fig. 2, the CPU unit 1 includes a CPU 11, a ROM 12, a RAM 13, and communication interfaces (hereinafter, referred to as "communication IF") 14 and 15. Additionally, the ROM 12 is one example of "the storing unit" of the present invention, and the communication IF 15 is one example of "the communication unit" of the present invention.

[0039] The CPU 11 has a function for executing the system program 121. The CPU 11 is a single core device and configured to execute a plurality of tasks in parallel with one another in a time sharing manner. That is, the CPU 11 executes multitask processes.

[0040] The ROM 12 is a nonvolatile memory such as a flash memory. In the ROM 12, the system program 121 is stored, and information L1 (see Fig. 4) relating to tasks to be executed in the CPU unit 1 is also stored. The information L1 relating to tasks includes names of the respective tasks, degrees of preference of the respective tasks, cycles in which the respective tasks are executed, configurations of the user programs to be executed in the respective tasks and a program organization unit (POU) for forming the user program for each of the tasks. Moreover, in the CPU unit 1, a control cycle having a predetermined time interval (for example, 1 ms) is adopted as a common cycle for the entire processes. In this case, with respect to the tasks, a task having a smaller value in the degree of preference is executed preferentially. Moreover, the user program is described in, for example, a ladder language.

[0041] In this case, each program organization unit is a minimum unit on program managements, and includes a program, a function, and a function block. Moreover, the function and the function block of the program organization unit can be re-used. That is, the function and the function block of the program organization unit can be commonly used in a plurality of user programs. Here, the function is a function for outputting a constant output value corresponding to the predetermined input value when a predetermined input value is inputted, and the function block is a function whose output value in response to the input value is not constant, because it is allowed to maintain an inner state.

[0042] In the example of Fig. 4, a task A is more preferentially executed than tasks B and C, and is executed in a cycle of 1 ms. That is, the task A uses a control cycle of one time (time interval of 1 ms) as its execution cycle. Moreover, the task A includes a user program UPa configured by programs Pa1 and Pa2. Moreover, the task A includes a function block FB1 to be utilized in the program Pa1. Additionally, the task A is one example of "the ref-

erence task" of the present invention.

[0043] The task B is more preferentially executed than the task C, and is executed in a cycle of 2 ms. That is, the task B uses a control cycle of two times (time interval of 2 ms) as its execution cycle. Moreover, the cycle of the task B is set to an integral multiple (2 times) of the cycle of the task A. Furthermore, the task B includes a user program UPb configured by programs Pb1 and Pb2. The task B includes a function block FB1 to be utilized in the program Pb2. Additionally, the task B is one example of "a task other than the reference task" of the present invention.

[0044] The task C is executed more subordinately than the tasks A and B, and is executed in a cycle of 4 ms. That is, the task C uses a control cycle of four times (time interval of 4 ms) as its execution cycle. Furthermore, the cycle of the task C is set to an integral multiple (4 times) of the cycle of the task A. The task C includes a user program UPc configured by a program Pc1. Further, the task C includes function blocks FB2 and FB3 to be utilized in the program Pc1. Additionally, the task C is one example of "a task other than the reference task" of the present invention.

[0045] The user programs UPa to UPc are stored in the ROM 12. For this reason, in the ROM 12, the programs Pa1, Pa2, Pb1, Pb2 and Pc1, each forming a program organization unit, and function blocks FB1 to FB3, each forming a program organization unit, are stored.

[0046] The RAM 13 is a volatile memory such as a DRAM. The RAM 13 has a function for temporarily storing a system program 121 or the like to be executed by the CPU 11. The communication IF 14 is connected to the internal bus 6, and is installed so as to allow the CPU unit 1 to communicate with the respective units via the internal bus 6. The communication IF 15 is connected to a tool device 80, and is installed so as to allow the CPU unit 1 to communicate with the tool device 80.

[0047] The tool device 80 is, for example, a personal computer, and has functions for forming and editing a user program to be executed in the PAC 100. Moreover, the tool device 80 has a function for downloading the user program from the PAC 100, and also has a function for uploading the user program to the PAC 100. Furthermore, the tool device 80 has a function for transmitting a switching instruction of operating modes of the CPU unit 1.

[0048] As shown in Fig. 3, the CPU unit 1 includes a task executing unit 21, a mode switching unit 22, a receiving unit 23 and a management unit 24. In this case, the task executing unit 21, the mode switching unit 22, the receiving unit 23 and the management unit 24 are realized when the CPU 11 (see Fig. 2) executes the system program 121. The task executing unit 21 is one example of "the task executing unit" of the present invention, and the mode switching unit 22 is one example of "the mode switching unit" of the present invention. Moreover, the receiving unit 23 is one example of "the receiving unit" of the present invention, and the management

unit 24 is one example of "the management unit" of the present invention.

[0049] The task executing unit 21 is configured so as to execute a plurality of tasks in parallel with one another and also execute the respective tasks in cycles based on the respective tasks. More specifically, the task executing unit 21 executes the tasks A, B and C in parallel with one another in a time sharing manner. Moreover, the task executing unit 21 executes the task A in a cycle of 1 ms, executes the task B in a cycle of 2 ms, and also executes the task C in a cycle of 4 ms.

[0050] Moreover, the task executing unit 21 is configured so as to execute a user program in each task at the time of an operation mode, and so as not to execute a user program in each task at the time of a program mode. More specifically, at the time of the operation mode, the task executing unit 21 executes the user program UPa and a program SPa other than the user program in the task A, executes the user program UPb and a program SPb other than the user program in the task B, and executes the user program UPc and a program SPc other than the user program in the task C. Moreover, at the time of the program mode, the task executing unit 21 executes only the program SPa other than the user program in the task A, executes only the program SPb other than the user program in the task B, and executes only the program SPc other than the user program in the task C. Additionally, the programs SPa, SPb and SPc other than the user program are programs forming one portion of the system program 121.

[0051] In this case, the operation mode is a mode for actually controlling the control object apparatus 60 (see Fig. 1) such as an industrial machine by executing the user program, and the program mode is a mode for allowing the user to edit a user program by using the tool device 80 (see Fig. 2) without executing the user program. That is, the operation mode is a mode in which control on the control object apparatus 60 is carried out, while the program mode is a mode in which no control on the control object apparatus 60 is carried out. For this reason, when the mode is switched from the operation mode to the program mode, the execution of the user program is stopped, and when the mode is switched from the program mode to the operation mode, the execution of the user program is started. In this case, the operation mode is one example of "the first mode" of the present invention, and the program mode is one example of "the second mode" of the present invention.

[0052] The mode switching unit 22 is configured to switch operating modes when a plurality of tasks executed by the task executing unit 21 are synchronized with one another (when execution cycles of the plurality of tasks are coincident with one another). Additionally, the operating modes include the operation mode and the program mode. More specifically, when the task A to be executed in the cycle of 1 ms, the task B to be executed in the cycle of 2 ms and the task C to be executed in the cycle of 4 ms are synchronized with one another, the

mode switching unit 22 switches the operating modes. In this case, the time at which the tasks are synchronized with one another refers to a point of time after completion of executions of a plurality of tasks carried out in cycles based on the respective tasks and prior to the simultaneous starts of the execution cycles of the plurality of tasks.

[0053] The receiving unit 23 has a function for receiving a switching instruction of operating modes. More specifically, the receiving unit 23 receives a switching instruction transmitted from the tool device 80 through the communication IF 15 (see Fig. 2). In this case, the switching instruction to be transmitted from the tool device 80 includes a switching instruction from the operation mode to the program mode (user program stopping instruction) and a switching instruction from the program mode to the operation mode (user program starting instruction).

[0054] The management unit 24, which is a scheduler, manages executions of a plurality of tasks by the task executing unit 21, and has a function for managing a timing in which the mode switching unit 22 switches the operating modes. More specifically, the management unit 24 is configured such that, when a plurality of tasks executed by the task executing unit 21 are synchronized with one another after the receiving unit 23 receives a switching instruction, the operating modes are switched by the mode switching unit 22.

[0055] Additionally, the CPU unit 1 is configured such that, even when multitask processes are executed at the time of the operation mode, an editing process of a user program can be carried out. That is, the CPU unit 1 is configured so as to carry out an on-line editing process.

[0056] Next, referring to Fig. 5, the following description will be given of operations, at the time of the operation mode, of the CPU unit 1 in accordance with the first embodiment. In this case, in the operation mode, user programs for the respective tasks are executed. Moreover, the following operations are carried out when the system program 121 is executed by the CPU 11 of the CPU unit 1 (see Fig. 2). That is, the system program 121 is a system program for allowing the CPU unit 1 to execute the following operations.

[0057] At the time of the operation mode, upon starting a control cycle, execution cycles of the tasks A, B and C are simultaneously started by the management unit 24 (see Fig. 3). Additionally, the task A is executed in a cycle of 1 ms, the task B is executed in a cycle of 2 ms, and the task C is executed in a cycle of 4 ms.

[0058] More specifically, first, the process of the task A having the highest degree of preference is started by the task executing unit 21 (see Fig. 3). At this time, since the tasks B and C are lower in the degree of preference than the task A, actual processes for the tasks B and C are not started and the tasks B and C are kept in stand-by states.

[0059] Then, the user program UPa of the task A and the program SPa other than the user program are executed, and upon completion of the process of the task A,

a process of the task B having a higher degree of preference than that of the task C is started. At this time, the task C is kept in the stand-by state. Then, the user program UPb of the task B is executed.

[0060] Next, after a lapse of 1 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the user program UPb of the task B is suspended by the management unit 24, and the execution of the task A is started. At this time, the tasks B and C are kept in stand-by states.

[0061] Moreover, the user program UPa of the task A and the program SPa other than the user program are executed, and upon completion of the process of the task A, the process of the suspended task B is resumed. Thus, the rest of the user program UPb of the task B and the program SPb other than the user program are executed. Upon completion of the process of the task B, the process of the task C is started so that the user program UPc of the task C is executed.

[0062] Next, after a lapse of 2 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the user program UPc of the task C is suspended by the management unit 24, and the execution of the task A is started. At this time, since the execution cycle of the task B has also been elapsed, the tasks B and C are kept in stand-by states. That is, after a lapse of 2 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, the execution cycles of the tasks A and B are again started simultaneously.

[0063] Then, the user program UPa of the task A and the program SPa other than the user program are executed, and upon completion of the task A, the process of the task B is started. At this time, the task C is kept in the stand-by state. Then, the user program UPb of the task B is executed.

[0064] Next, after a lapse of 3 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the user program UPb of the task B is suspended by the management unit 24, and the execution of the task A is started. At this time, the tasks B and C are kept in stand-by states.

[0065] Moreover, the user program UPa of the task A and the program SPa other than the user program are executed, and upon completion of the process of the task A, the process of the suspended task B is resumed. Thus, the rest of the user program UPb of the task B and the program SPb other than the user program are executed. Moreover, upon completion of the process of the task B, the process of the suspended task C is resumed. Thus, the rest of the user program UPc of the task C and the program SPc other than the user program are executed. Then, after a lapse of 4 ms from the simultaneous starts of the execution cycles of the tasks A, B and C after the completion of the process of the task C, since the execution cycle of the task A has been elapsed, the execution

of the task A by the task executing unit 21 is carried out by the management unit 24. At this time, since the execution cycles of the tasks B and C have also been elapsed, the tasks B and C are kept in stand-by states. That is, after a lapse of 4 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, the execution cycles of the tasks A, B and C are again started simultaneously. Thereafter, the above-mentioned operations are repeatedly carried out.

[0066] In this manner, in the operation mode, the user program UPa of the task A, the user program UPb of the task B and the user program UPc of the task C are executed so that the control object apparatus 60 (see Fig. 1) is controlled by the PAC 100 (see Fig. 1).

[0067] Next, referring to Fig. 6, the following description will be given of operations, at the time of the program mode, of the CPU unit 1 in accordance with the first embodiment. In the program mode, only the programs other than the user programs are executed without executing the user programs of the respective tasks. Moreover, the following operations are carried out when the system program 121 is executed by the CPU 11 of the CPU unit 1 (see Fig. 2). That is, the system program 121 is a system program for allowing the CPU unit 1 to execute the following operations.

[0068] In the program mode, upon starting a control cycle, execution cycles of the tasks A, B and C are simultaneously started by the management unit 24 (see Fig. 3). Additionally, the task A is executed in a cycle of 1 ms, the task B is executed in a cycle of 2 ms, and the task C is executed in a cycle of 4 ms.

[0069] More specifically, first, the process of the task A having the highest degree of preference is started by the task executing unit 21 (see Fig. 3). At this time, since the tasks B and C are lower in the degree of preference than the task A, actual processes for these are not started and these are kept in stand-by states.

[0070] Moreover, the program SPa other than the user program of the task A is executed, and upon completion of the process of the task A, a process of the task B having a higher degree of preference than that of the task C is started. At this time, the task C is kept in the stand-by state.

[0071] Then, the program SPb other than the user program of the task B is executed, and upon completion of the process of the task B, a process of the task C is started. Thereafter, the program SPc other than the user program of the task C is executed, and the process of the task C is completed.

[0072] Next, after a lapse of 1 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the execution of the task A by the task executing unit 21 is carried out by the management unit 24. Thus, the program SPa other than the user program of the task A is executed so that the process of the task A is completed.

[0073] Next, after a lapse of 2 ms from the simultane-

ous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the execution of the task A by the task executing unit 21 is carried out by the management unit 24. Thus, the process of the task A is started. At this time, since the execution cycle of the task B has also been elapsed, the task B is kept in the stand-by state. That is, after a lapse of 2 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, the execution cycles of the tasks A and B are again started simultaneously.

[0074] Then, the program SPa other than the user program of the task A is executed, and upon completion of the task A, the process of the task B is started. Moreover, the program SPb other than the user program of the task B is executed, and the process of the task B is completed.

[0075] Next, after a lapse of 3 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the execution of the task A by the task executing unit 21 is carried out by the management unit 24. Thus, the program SPa other than the user program of the task A is executed so that the process of the task A is completed.

[0076] Next, after a lapse of 4 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, since the execution cycle of the task A has been elapsed, the execution of the task A by the task executing unit 21 is carried out by the management unit 24. At this time, since the execution cycles of the tasks B and C have also been elapsed, the tasks B and C are kept in stand-by states. That is, after a lapse of 4 ms from the simultaneous starts of the execution cycles of the tasks A, B and C, the execution cycles of the tasks A, B and C are again started simultaneously. Thereafter, the above-mentioned operations are carried out repeatedly.

[0077] In this manner, in the program mode, the user can edit the user programs UPa, UPb and UPc by using the tool device 80 (see Fig. 2) without executing the user program UPa of the task A, the user program UPb of the task B and the user program UPc of the task C. Additionally, in the program mode, executions of the user programs in the respective tasks are omitted, however, the respective tasks are executed in the same execution cycles of those of the operation mode.

[0078] Next, referring to Fig. 7, the following description will be given of operations, at the time of switching modes from the operation mode to the program mode, of the CPU unit 1 in accordance with the first embodiment. Additionally, the switching modes from the operation mode to the program mode means that the control of the control object apparatus 60 (see Fig. 1) by the PAC 100 (see Fig.

1) is stopped by stopping the execution of the user program. Moreover, the following description will be given of operations in the case that a switching instruction (stopping instruction) is transmitted from the tool device 80 (see Fig.

2) during the operation mode, and the receiving unit 23 (see Fig. 3) receives the switching instruction at point of time P1 by receiving the switching instruction by the communication IF 15 (see Fig. 2).

[0079] Moreover, the following operations are carried out when the system program 121 is executed by the CPU 11 of the CPU unit 1 (see Fig. 2). That is, the system program 121 is a system program for allowing the CPU unit 1 to execute the following operations.

[0080] In the operation mode, upon starting a control cycle, execution cycles of the tasks A, B and C are simultaneously started by the management unit 24 (see Fig. 3). In this case, the processes of the tasks A, B and C are the same as those as described earlier. Moreover, the switching instruction is received by the receiving unit 23 at the point of time P1.

[0081] For this reason, when the tasks A, B and C are next synchronized with one another (when the execution cycles of the tasks A, B and C are simultaneously started), the management unit 24 allows the mode switching unit 22 (see Fig. 3) to switch modes from the operation mode to the program mode. That is, at point of time P2, when the execution cycles of the tasks A, B and C are made coincident with one another, the operation mode is switched to the program mode.

[0082] Thus, even after the receiving unit 23 receives the switching instruction, up to the time when the execution cycles of the tasks A, B and C are simultaneously started at point of time P2, the user program UPa of the task A, the user program UPb of the task B and the user program UPc of the task C are executed. Therefore, the executions of the user program UPa of the task A, the user program UPb of the task B and the user program UPc of the task C are not interrupted in the middle thereof.

[0083] Moreover, when at the point of time P2, the execution cycles of the tasks A, B and C are simultaneously started, the tasks A, B and C are executed in the program mode. Thus, only the program SPa other than the user program is executed in the task A, only the program SPb other than the user program is executed in the task B, and only the program SPc other than the user program is executed in the task C.

[0084] Next, referring to Fig. 8, the following description will be given of operations, at the time of switching modes from the program mode to the operation mode, of the CPU unit 1 in accordance with the first embodiment. Additionally, the switching modes from the program mode to the operation mode means that the control of the control object apparatus 60 (see Fig. 1) by the PAC 100 (see Fig. 1) is started by starting the execution of the user program. Moreover, the following description will be given of operations in the case that a switching instruction (starting instruction) is transmitted from the tool device 80 (see Fig. 2) during the program mode, and the receiving unit 23 (see Fig. 3) receives the switching instruction at point of time P3 by receiving the switching instruction by the communication IF 15 (see Fig. 2).

[0085] Moreover, the following operations are carried out when the system program 121 is executed by the CPU 11 of the CPU unit 1 (see Fig. 2). That is, the system program 121 is a system program for allowing the CPU unit 1 to execute the following operations.

[0086] In the program mode, upon starting a control cycle, execution cycles of the tasks A, B and C are simultaneously started by the management unit 24 (see Fig. 3). In this case, the processes of the tasks A, B and C are the same as those as described above. Moreover, the switching instruction is received by the receiving unit 23 at the point of time P3.

[0087] For this reason, when the tasks A, B and C are next synchronized with one another (when the execution cycles of the tasks A, B and C are simultaneously started), the management unit 24 allows the mode switching unit 22 (see Fig. 3) to switch modes from the program mode to the operation mode. That is, when the execution cycles of the tasks A, B and C are made coincident with one another at point of time P4, the program mode is switched to the operation mode.

[0088] Thus, even after the receiving unit 23 receives the switching instruction, up to the time when the execution cycles of the tasks A, B and C are simultaneously started at the point of time P4, none of the user program UPa of the task A, the user program UPb of the task B and the user program UPc of the task C are executed. Therefore, for example, it is possible to prevent the occurrence of such a state in which the user program UPa of the task A and the user program UPb of the task B are executed, with the user program UPc of the task C being not executed.

[0089] Moreover, when the execution cycles of the tasks A, B and C are simultaneously started at the point of time P4, the tasks A, B and C are executed in the operation mode. Thus, in the task A, the user program UPa and the program SPa other than the user program are executed, in the task B, the user program UPb and the program SPb other than the user program are executed, and in task C, the user program UPc and the program SPc other than the user program are executed.

[0090] In the first embodiment, as described above, in the case that a plurality of tasks to be executed (tasks A, B and C) are synchronized with one another after the receiving unit 23 receives the switching instruction, switching operating modes can prevent the execution of the user program of each task from being suspended in the middle and prevent only the user program of one portion of the task from being executed. Therefore, while multitask processes are being executed, the switching of operating modes can be carried out appropriately. As a result, since correct operation results can be obtained in the user programs of the respective tasks, the PAC 100 is allowed to control the control object apparatus 60 normally.

[0091] Moreover, in the first embodiment, the CPU 11 is a single core device, and by executing a plurality of tasks in parallel with one another in a time sharing man-

ner, the processing capability of the CPU 11 can be efficiently utilized.

[0092] Furthermore, in the first embodiment, since the cycles of the tasks B and C are set to integral multiples of the cycle of the task A, the tasks A, B and C can be easily synchronized with one another.

<Second Embodiment>

[0093] First, referring to Figs. 9 to 10, the following description will be given of a configuration of a CPU unit 1 a in accordance with a second embodiment. Additionally, the same components as those of the CPU unit 1 of the first embodiment are indicated by the same reference numerals, and overlapped descriptions thereof will be omitted.

[0094] As shown in Fig. 9, the CPU unit 1 a in which a system program 121 a is installed is configured such that operations of the CPU unit 1 a are controlled by executing the system program 121 a. In this case, the CPU unit 1 a is one example of "the control device" of the present invention.

[0095] The CPU unit 1 a includes a CPU 11 a, a ROM 12a, a RAM 13, and communication IF 14 and 15. Additionally, the ROM 12a is one example of "the storing unit" of the present invention.

[0096] The CPU 11 a has a function for executing the system program 121 a. The CPU 11 a is a dual core device and configured to execute a plurality of tasks in parallel with one another simultaneously. That is, the CPU 11 a executes multitask processes. The other configurations of the CPU 11 a are the same as those of the aforementioned CPU 11.

[0097] In the ROM 12a, the system program 121 a is stored, and information L2 (see Fig. 10) relating to tasks to be executed in the CPU unit 1 a is also stored. The information L2 relating to tasks includes names of the respective tasks, degrees of preference of the respective tasks, cycles in which the respective tasks are executed, configurations of the user programs to be executed in the respective tasks and a program organization unit for forming the user program for each of the tasks. Additionally, the other configurations of the ROM 12a are the same as those of the aforementioned ROM 12.

[0098] In the example of Fig. 10, a task D is more preferentially executed than a task E, and is executed in a cycle of 1 ms. That is, the task D uses a control cycle of one time (time interval of 1 ms) as its execution cycle. Moreover, the task D includes a user program UPd configured by programs Pd1 and Pd2. Moreover, the task D includes a function block FB4 to be utilized in the program Pd1. Additionally, the task D is one example of "the reference task" of the present invention.

[0099] Moreover, a task E is executed more subordinately than the task D, and is executed in a cycle of 2 ms. That is, the task E uses a control cycle of two times (time interval of 2 ms) as its execution cycle. Moreover, the cycle of the task E is set to an integral multiple (two

times) of the cycle of the task D. Furthermore, the task E includes a user program UPe configured by programs Pe1 and Pe2. The task E includes a function block FB4 to be utilized in the program Pe2. Additionally, the task E is one example of "a task other than the reference task" of the present invention.

[0100] Next, referring to Fig. 11, the following description will be given of operations at the time of the operation mode in the CPU unit 1 a in accordance with the second embodiment. In this case, in the operation mode, a user program of each task and a program other than the user program are executed. Moreover, the following operations are carried out when the system program 121 a is executed by the CPU 11 a of the CPU unit 1 a (see Fig. 9). That is, the system program 121 a is a system program for allowing the CPU unit 1 a to execute the following operations.

[0101] At the time of the operation mode, upon starting a control cycle, execution cycles of the tasks D and E are simultaneously started by the management unit 24 (see Fig. 3). Additionally, the task D is executed in a cycle of 1 ms, and the task E is executed in a cycle of 2 ms.

[0102] In the CPU unit 1 a of the second embodiment, the task D and the task E are processed in parallel with each other. Moreover, after a lapse of 1 ms from the simultaneous starts of the execution cycles of the tasks D and E after the completion of the task D, since the execution cycle of the task D has been elapsed, the management unit 24 allows the task executing unit 21 to execute the task D. Thereafter, after a lapse of 2 ms from the simultaneous starts of the execution cycles of the tasks D and E after the completion of the process of the tasks D and E, since the execution cycles of the tasks D and E have been elapsed, the executions of the tasks D and E by the task executing unit 21 are again carried out by the management unit 24. Thereafter, the above-mentioned operations are repeatedly carried out.

[0103] Additionally, operations at the time of the program mode of the CPU unit 1 a of the second embodiment, operations at the time of switching modes from the operation mode to the program mode thereof, and operations at the time of switching modes from the program mode to the operation mode thereof are the same as those of the CPU unit 1 of the first embodiment except that the tasks D and E are simultaneously executed in parallel with each other.

[0104] As described above, in the second embodiment, since the CPU 11 a is a dual core device, and executes a plurality of tasks simultaneously in parallel with one another so that it is possible to improve the task processing capability.

[0105] In this case, the other effects of the second embodiment are the same as those of the first embodiment.

[0106] Additionally, the present embodiments disclosed above are exemplary only in all the aspects, and do not form the basis of a limited interpretation. Therefore, the technical scope of the present invention is not interpreted only by the above-mentioned embodiments,

but defined based upon the description of the claims. Moreover, the technical scope of the present invention includes all modifications made within the meaning and scope equal to the claims.

[0107] For example, the first embodiment has exemplified a configuration in which the PAC 100 is provided with the CPU unit 1, the input unit 2, the output unit 3, communication unit 4 and the power supply unit 5; however, not limited to this configuration, the PAC may have another unit, and none of the input unit and the communication unit may be installed in the PAC.

[0108] Moreover, the first embodiment has exemplified a configuration in which the PAC 100 is constituted by a plurality of units; however, not limited by this configuration, the PAC may be formed by one casing, with functions of the respective units being incorporated in the casing.

[0109] Furthermore, the first embodiment has exemplified a configuration in which a switching instruction transmitted from the tool device 80 is received by the receiving unit 23; however, not limited by this configuration, a switching instruction transmitted by another PAC 70 via a communication unit 4 may be received by the receiving unit 23. Moreover, a switching instruction generated by a user program executed by the CPU unit 1 may be received by the receiving unit 23.

[0110] The first embodiment has exemplified a configuration in which, when the CPU 11 executes the system program 121, the task executing unit 21, the mode switching unit 22, the receiving unit 23 and management unit 24 are realized; however, not limited to this configuration, the task executing unit 21, the mode switching unit 22, the receiving unit 23 and the management unit 24 may be respectively configured by hardware units.

[0111] Furthermore, with respect to tasks in the first and second embodiments, in addition to tasks that are regularly executed (fixed-cycle tasks), a task (event task) that is executed when a preliminary set execution condition is satisfied may be included.

[0112] Moreover, the second embodiment has exemplified a configuration in which the CPU 11 is a dual core device; however, not limited to this configuration, for example, the CPU may be a quad core device. That is, the number of cores of the CPU may be any number.

DESCRIPTION OF SYMBOLS

[0113]

1, 1a	CPU unit (control device)
12, 12a	ROM (storing unit)
15	communication IF (communication unit)
21	task executing unit (task execution means)
22	mode switching unit (mode switching means)
23	receiving unit (receiving means)
24	management unit (management means)
60	control object apparatus

80	tool device
121, 121	a system program

5 Claims

1. A control device (1, 1 a) for controlling a control object apparatus (60) comprising:

10 a task executing unit (21) configured to execute a plurality of tasks in parallel with one another and to execute each task in a cycle based on the task,
 15 a mode switching unit (22) configured to switch operation modes of the task executing unit (21), a receiving unit (23) configured to receive a switching instruction for operating modes in the task executing unit (21),
 20 a storing unit (12, 12a) configured to store a user program; and
 25 a communication unit (15) configured to communicate with a tool device (80) for editing the user program stored in the storing unit (12, 12a), wherein the task executing unit (21) is configured such that in a first mode, a user program is executed for each task, and in a second mode, a user program is not executed for each task, wherein the mode switching unit (22) is configured such that upon receipt of a switching instruction for operating modes of the task executing unit (21), the mode switching unit (22) does not switch the operating modes of the task executing unit (21) until the plurality of tasks to be executed by the task executing unit (21) have been synchronized with each other, and the operating mode of the task executing unit (21) is switched when the plurality of tasks are synchronized with each other,
 30 wherein after the receipt of the switching instruction by the receiving unit (23), when the plurality of tasks which are being executed by the task executing unit (21) are synchronized with one another, the mode switching unit (22) switches the operating mode in the task executing unit (21), and
 35 wherein the receiving unit (23) receives the switching instruction transmitted from the tool device (80) via the communication unit (15).

- 50 2. The control device (1, 1 a) according to claim 1, wherein the task executing unit (21) is configured so as to execute a plurality of tasks simultaneously in parallel with one another.
- 55 3. The control device (1, 1 a) according to claim 1, wherein the task executing unit (21) is configured so as to execute a plurality of tasks in parallel with one another in a time sharing manner.

4. The control device (1, 1a) according to claim 1, wherein the plurality of tasks comprise a reference task and tasks other than the reference task, with each of the tasks other than the reference task being set to have a cycle having an integral multiple of the cycle of the reference task. 5
5. A system program, which is executed by a control device (1, 1 a) comprising a first mode for controlling a control object apparatus (60) by executing a user program in each task and a second mode in which no user program is executed in each task, the system program being configured to allow the control device (1, 1 a) to execute: 10
- a sequence for executing a plurality of tasks in parallel in a cycle based on each task, 15
- a sequence in which upon receipt of a switching instruction for operating modes between the first mode and the second mode, the operating modes are not switched until the plurality of tasks to be executed have been synchronized with each other and the operating mode of the task executing unit (21) is switched between the first mode and the second mode, when the plurality of tasks executed are synchronized with each other, 20
- a sequence for receive a switching instruction for operating modes, and, after the receipt of the switching instruction, switching the operating mode when the plurality of tasks which are being executed are synchronized with one another, 25
- a sequence for communicating with a tool device (80) for editing a user program stored in a storing unit (12, 12a), 30
- wherein the switching instruction transmitted from the tool device (80) is received via the communication unit (15). 35
6. A recording media for recording the system program according to claim 5, allowing a computer to read the system program. 40

Patentansprüche 45

1. Steuereinrichtung (1, 1a) zum Steuern einer Steuerelementvorrichtung (60) mit:
- einer Task-Ausführungseinheit (21), die dazu eingerichtet ist, mehrere Tasks parallel zueinander auszuführen und einen jeden Task in einem Zyklus abhängig von dem Task auszuführen, 50
- einer Moduswechseleinheit (22), die dazu eingerichtet ist, Betriebsmodi der Task-Ausführungseinheit (21) zu wechseln, 55
- einer Empfangseinheit (23), die dazu eingerich-

tet ist, eine Wechsellanweisung für Betriebsmodi in der Task-Ausführungseinheit (21) zu empfangen,

einer Speichereinheit (12, 12a), die dazu eingerichtet ist, ein Benutzerprogramm zu speichern; und

einer Kommunikationseinheit (15), die dazu eingerichtet ist, mit einer Werkzeugeinrichtung (80) zum Bearbeiten des in der Speichereinheit (12, 12a) gespeicherten Benutzerprogramms zu kommunizieren,

wobei die Task-Ausführungseinheit (21) dergestalt eingerichtet ist, dass in einem ersten Modus für einen jeden Task ein Benutzerprogramm ausgeführt wird und in einem zweiten Modus für einen jeden Task ein Benutzerprogramm nicht ausgeführt wird,

wobei die Moduswechseleinheit (22) dergestalt eingerichtet ist, dass die Moduswechseleinheit (22) bei Empfang einer Wechsellanweisung bezüglich Betriebsmodi der Task-Ausführungseinheit (21) die Betriebsmodi der Task-Ausführungseinheit (21) so lange nicht wechselt, bis die von der Task-Ausführungseinheit (21) auszuführenden mehreren Tasks miteinander synchronisiert worden sind, und der Betriebsmodus der Task-Ausführungseinheit (21) gewechselt wird, wenn die mehreren Tasks miteinander synchronisiert sind,

wobei nach Empfang der Wechsellanweisung durch die Empfangseinheit (23), wenn die mehreren von der Task-Ausführungseinheit (21) ausgeführten Tasks miteinander synchronisiert sind, die Moduswechseleinheit (22) den Betriebsmodus in der Task-Ausführungseinheit (21) wechselt und

wobei die Empfangseinheit (23) die von der Werkzeugeinrichtung (80) übertragene Wechsellanweisung über die Kommunikationseinheit (15) empfängt.

2. Steuereinrichtung (1, 1a) nach Anspruch 1, wobei die Task-Ausführungseinheit (21) dazu eingerichtet ist, mehrere Tasks gleichzeitig parallel zueinander auszuführen.
3. Steuereinrichtung (1, 1a) nach Anspruch 1, wobei die Task-Ausführungseinheit (21) dazu eingerichtet ist, mehrere Tasks auf eine Time-Sharing-Weise parallel zueinander auszuführen.
4. Steuereinrichtung (1, 1a) nach Anspruch 1, wobei die mehreren Tasks einen Referenz-Task und andere Tasks als den Referenz-Task umfassen, wobei ein jeder der anderen Tasks als der Referenz-Task mit einem Zyklus eingerichtet ist, der ein ganzzahliges Vielfaches des Zyklus des Referenz-Tasks ist.

5. Systemprogramm, das von einer Steuereinrichtung (1, 1a) ausgeführt wird und einen ersten Modus zum Steuern einer Steuerobjektvorrichtung (60) durch Ausführen eines Benutzerprogramms in einem jeden Task und einen zweiten Modus, in dem in einem jeden Task kein Benutzerprogramm ausgeführt wird, umfasst, wobei das Systemprogramm dazu eingerichtet ist, die Steuereinrichtung (1, 1a) ausführen zu lassen:

eine Sequenz zum Ausführen mehrerer Tasks parallel in einem Zyklus abhängig von einem jeden Task,
 eine Sequenz, bei welcher bei Empfang einer Wechsellanweisung bezüglich Betriebsmodi zwischen dem ersten Modus und dem zweiten Modus die Betriebsmodi so lange nicht gewechselt werden, bis die auszuführenden mehreren Tasks miteinander synchronisiert worden sind, und der Betriebsmodus der Task-Ausföhrungseinheit (21) zwischen dem ersten Modus und dem zweiten Modus gewechselt wird, wenn die ausgeführten mehreren Tasks miteinander synchronisiert sind,
 eine Sequenz zum Empfangen einer Wechsellanweisung bezüglich Betriebsmodi und, nach dem Empfang der Wechsellanweisung, Wechseln des Betriebsmodus, wenn die ausgeführten mehreren Tasks miteinander synchronisiert sind,
 eine Sequenz zum Kommunizieren mit einer Werkzeugeinrichtung (80) zum Bearbeiten eines in einer Speichereinheit (12, 12a) gespeicherten Benutzerprogramms,
 wobei die von der Werkzeugeinrichtung (80) übertragene Wechsellanweisung über die Kommunikationseinheit (15) empfangen wird.

6. Aufzeichnungsmedium zum Aufzeichnen des Systemprogramms nach Anspruch 5, was einem Computer erlaubt, das Systemprogramm zu lesen.

Revendications

1. Dispositif de commande (1, 1a) destiné à commander un appareil formant objet de commande (60) comprenant :

une unité d'exécution de tâches (21) conçue pour exécuter une pluralité de tâches en parallèle les unes des autres et pour exécuter chaque tâche dans un cycle en fonction de la tâche,
 une unité de commutation de modes (22) conçue pour commuter des modes de fonctionnement de l'unité d'exécution de tâches (21),
 une unité de réception (23) conçue pour recevoir une instruction de commutation de modes de

fonctionnement dans l'unité d'exécution de tâches (21),
 une unité de stockage (12, 12a) conçue pour stocker un programme utilisateur ; et
 une unité de communication (15) conçue pour communiquer avec un dispositif formant outil (80) destiné à éditer le programme utilisateur stocké dans l'unité de stockage (12, 12a), dans lequel l'unité d'exécution de tâches (21) est conçue de manière que dans un premier mode, un programme utilisateur soit exécuté pour chaque tâche, et dans un second mode, un programme utilisateur ne soit pas exécuté pour chaque tâche,
 dans lequel l'unité de commutation de modes (22) est conçue de manière qu'à la réception d'une instruction de commutation de modes de fonctionnement de l'unité d'exécution de tâches (21), l'unité de commutation de modes (22) ne commute pas les modes de fonctionnement de l'unité d'exécution de tâches (21) tant que les tâches de la pluralité de tâches devant être exécutées par l'unité d'exécution de tâches (21) n'ont pas été synchronisées les unes avec les autres, et que le mode de fonctionnement de l'unité d'exécution de tâches (21) soit commuté lorsque les tâches de la pluralité de tâches sont synchronisées les unes avec les autres,
 dans lequel après la réception de l'instruction de commutation par l'unité de réception (23), lorsque les tâches de la pluralité de tâches qui sont en cours d'exécution par l'unité d'exécution de tâches (21) sont synchronisées les unes avec les autres, l'unité de commutation de modes (22) commute le mode de fonctionnement dans l'unité d'exécution de tâches (21), et
 dans lequel l'unité de réception (23) reçoit l'instruction de commutation émise depuis le dispositif formant outil (80) par l'intermédiaire de l'unité de communication (15).

2. Dispositif de commande (1, 1a) selon la revendication 1, dans lequel l'unité d'exécution de tâches (21) est conçue de manière à exécuter une pluralité de tâches simultanément, en parallèle les unes des autres.
3. Dispositif de commande (1, 1a) selon la revendication 1, dans lequel l'unité d'exécution de tâches (21) est conçue de manière à exécuter une pluralité de tâches en parallèle les unes des autres, en temps partagé.
4. Dispositif de commande (1, 1a) selon la revendication 1, dans lequel la pluralité de tâches comprend une tâche de référence et des tâches autres que la tâche de référence, chacune des tâches autres que la tâche de référence étant définie comme ayant un

cycle représentant un multiple entier du cycle de la tâche de référence.

5. Programme système qui est exécuté par un dispositif de commande (1, 1a) comprenant un premier mode destiné à commander un appareil formant objet de commande (60) en exécutant un programme utilisateur dans chaque tâche, et un second mode dans lequel aucun programme utilisateur n'est exécuté dans chaque tâche, le programme système étant conçu pour permettre au dispositif de commande (1, 1a) d'exécuter :

une séquence destinée à exécuter une pluralité de tâches en parallèle, dans un cycle en fonction de chaque tâche, 5
 une séquence dans laquelle à la réception d'une instruction de commutation des modes de fonctionnement entre le premier mode et le second mode, les modes de fonctionnement ne sont pas commutés tant que les tâches de la pluralité de tâches à exécuter n'ont pas été synchronisées les unes avec les autres, et le mode de fonctionnement de l'unité d'exécution de tâches (21) est commuté entre le premier mode et le second mode, lorsque les tâches de la pluralité de tâches exécutées sont synchronisées les unes avec les autres, 10
 une séquence destinée à recevoir une instruction de commutation des modes de fonctionnement, et après la réception de l'instruction de commutation, à commuter le mode de fonctionnement, lorsque les tâches de la pluralité de tâches qui sont en cours d'exécution sont synchronisées les unes avec les autres, 15
 une séquence destinée à communiquer avec un dispositif formant outil (80) destiné à éditer un programme utilisateur stocké dans une unité de stockage (12, 12a), 20
 dans lequel l'instruction de commutation émise depuis le dispositif formant outil (80) est reçue par l'intermédiaire de l'unité de communication (15). 25
 30
 35
 40

6. Support d'enregistrement destiné à enregistrer le programme système selon la revendication 5, permettant à un ordinateur de lire le programme système. 45

50

55

Fig. 1

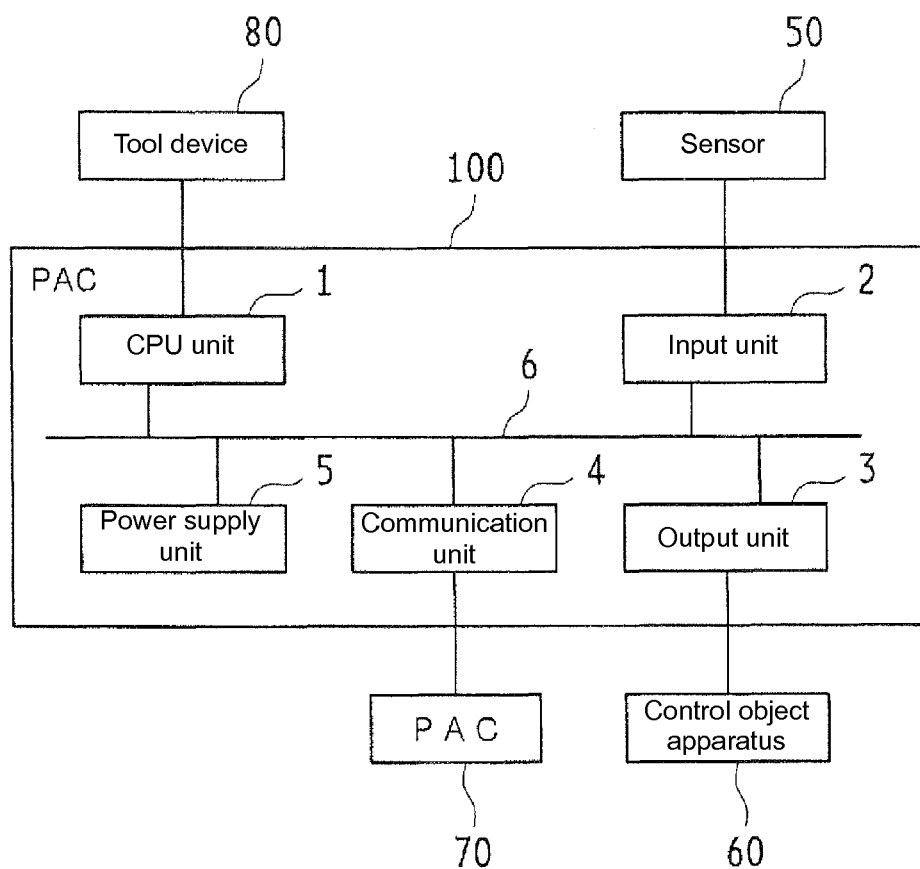


Fig. 2

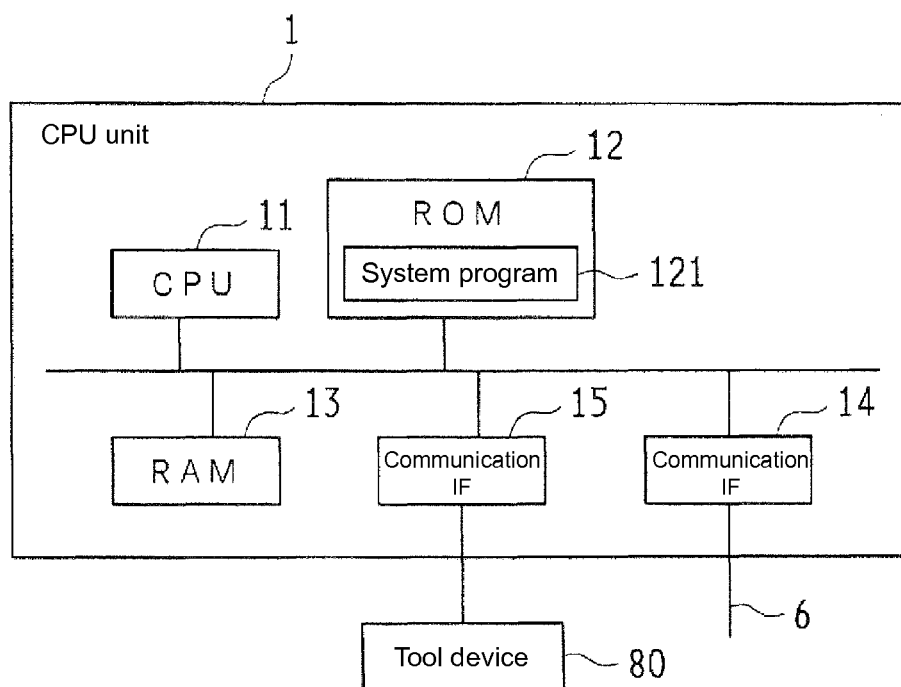


Fig. 3

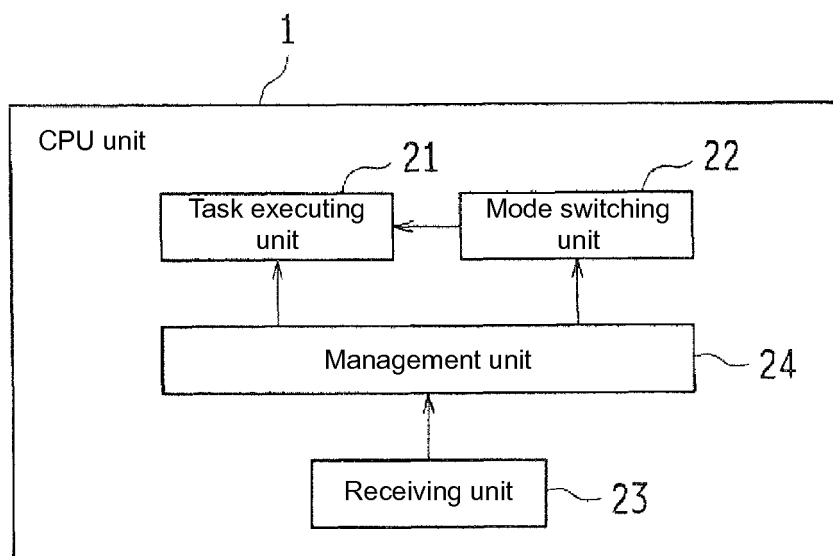


Fig. 4

L1

Name	Degree of preference	Cycle	User program	Program organization unit of user program	
				Program	Function-Function block
Task A	1	1ms	UPa	Pa1, Pa2	FB1
Task B	2	2ms	UPb	Pb1, Pb2	FB1
Task C	3	4ms	UPc	Pc1	FB2, FB3

Fig. 5

At the time of operation mode

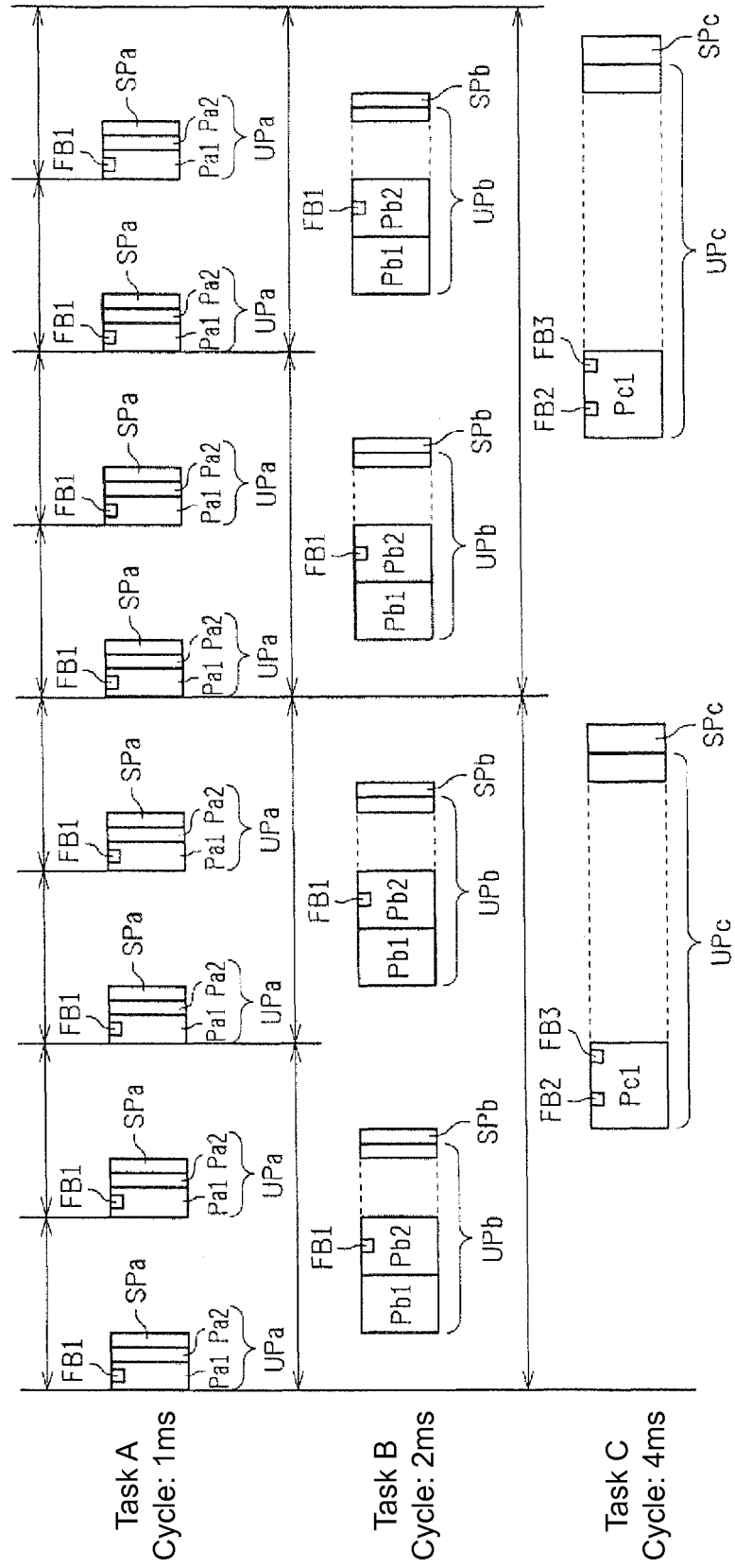


Fig. 6

At the time of program mode

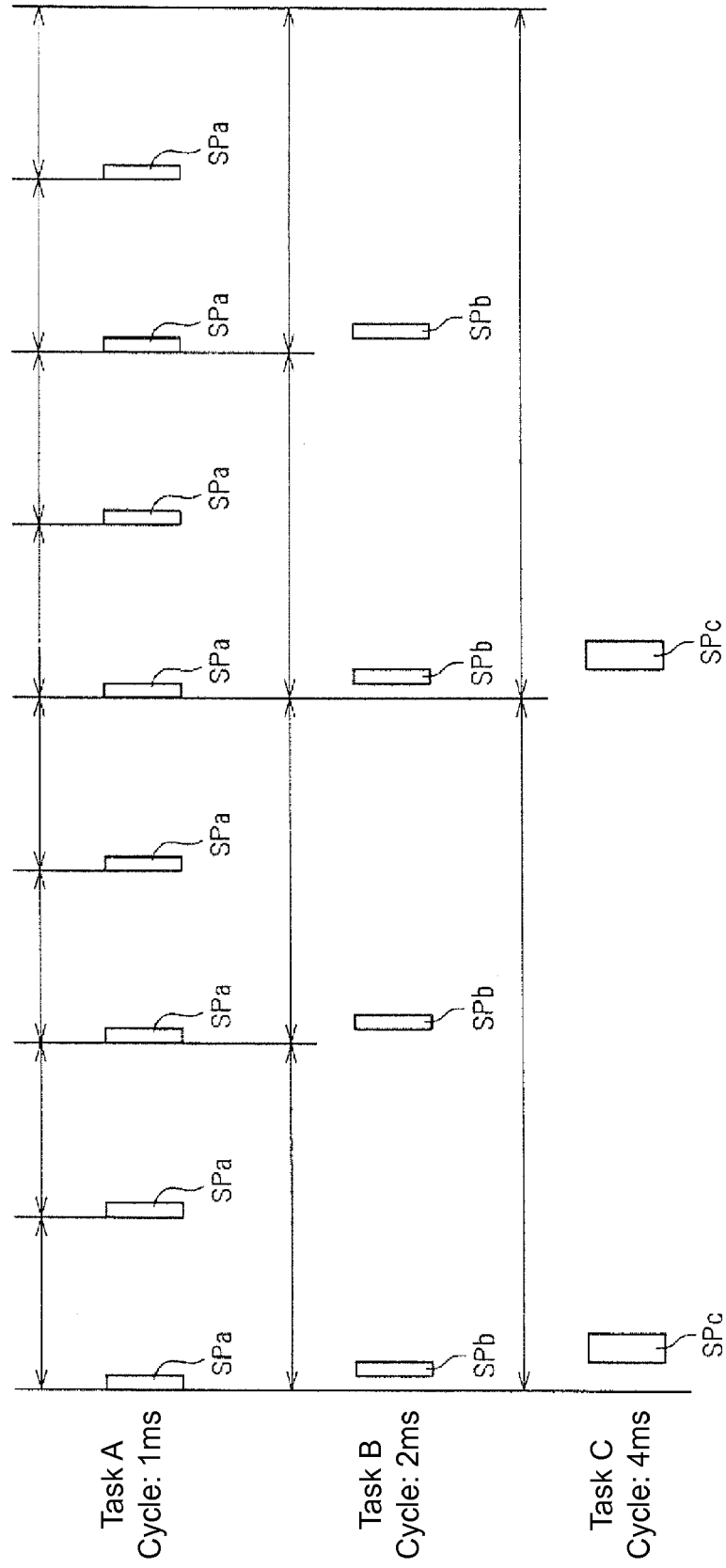


Fig. 7

Upon switching from operation mode to program mode

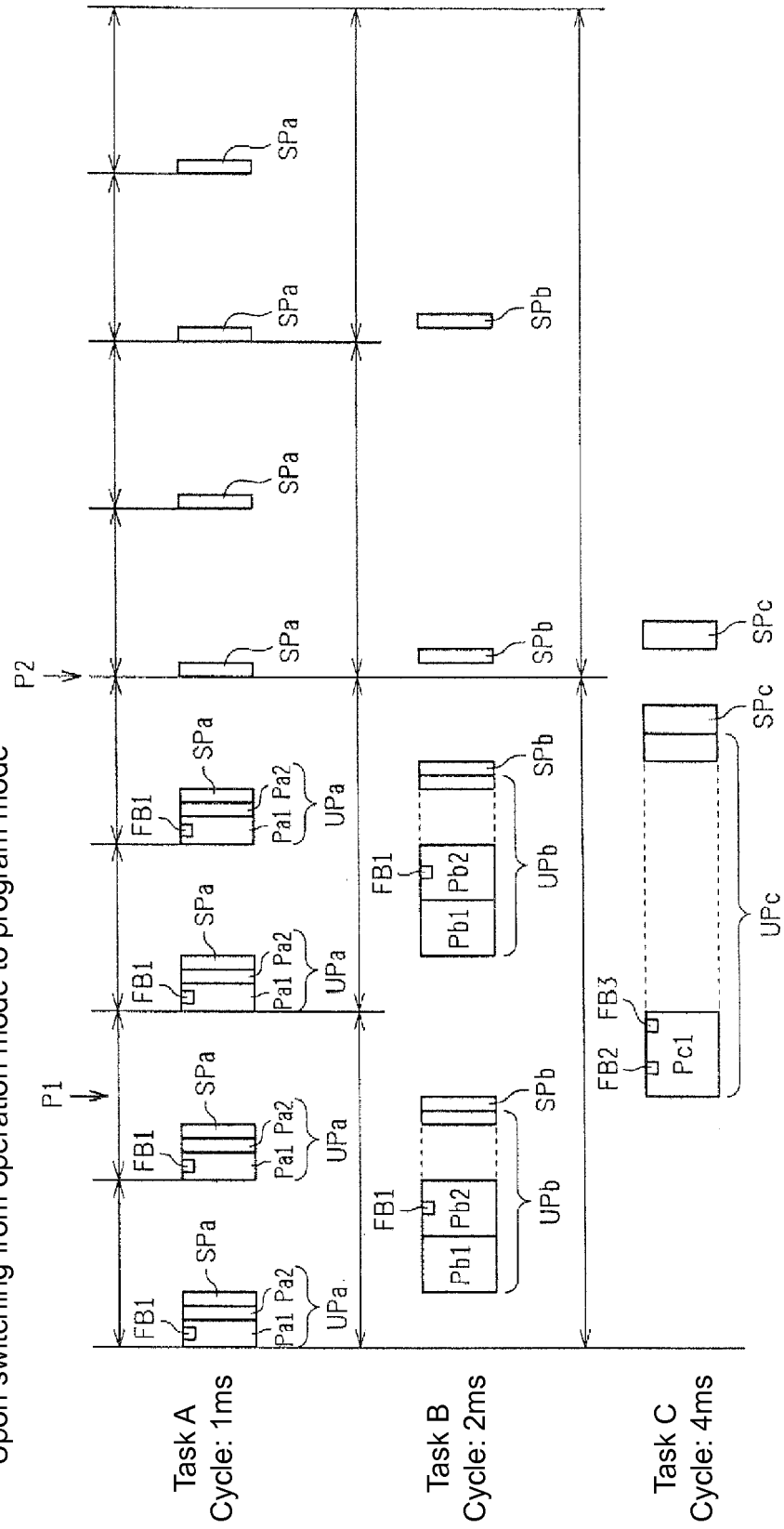


Fig. 9

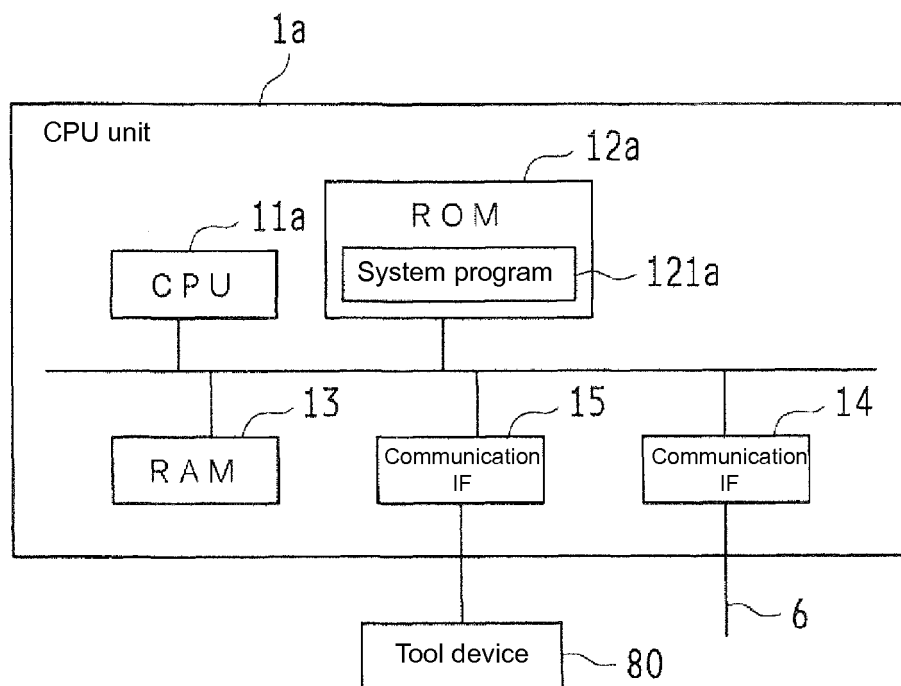


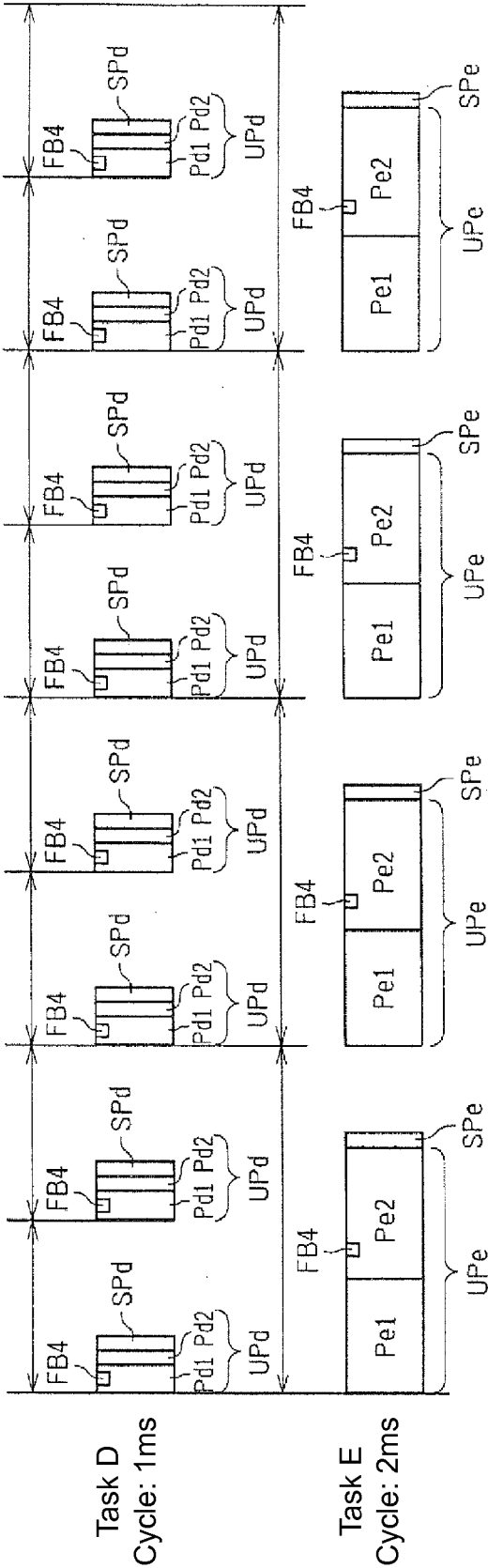
Fig. 10

L2

Name	Degree of preference	Cycle	User program	Program organization unit of user program	
				Program	Function-Function block
Task D	1	1ms	UPd	Pd1, Pd2	FB4
Task E	2	2ms	UPe	Pe1, Pe2	FB4

Fig. 11

At the time of operation mode



REFERENCES CITED IN THE DESCRIPTION

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